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E-voting with Receipt-freeness and Universal Verifiability Using E-voting Sheet

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1 Introduction

Sako and Killian[SK95] proposed mix-net e-voting system which satisfies receipt-freeness and universal verifiability in Eurocrypt'95. Michels and Horster pointed out that Sako-Killian scheme does not satisfy robustness and privacy[MH96]. Golle et al.[GJS04] proposed universal re-encryption mix net which satisfies correctness and communication privacy. In this paper, we propose a novel mix-net e-voting system using Golle et al.'s universal re-encryption mix-net which satisfies receipt-free and universal verifiability as well as robustness and privacy. To achieve universal verifiability, receipt-freeness, and privacy, we introduce firstly *E-voting Sheet* which only a voter with *E-voting Sheet* can cast a vote and an *Overwritable Bulletin Board* which can be overwritten contents of bulletin board by each mix-center. To achieve anonymity of encrypted voting content and robustness, we use Golle et al.'s universal re-encryption mix-net.

2 E-voting Procedure

2.1 Entities

Voter V_i ($\{i \mid i = 1, \dots, z\}$): A voter cast a vote only by an election rule.

Mix center C_i ($\{j \mid j = 1, \dots, n\}$)

- Each mix-center generates a random encryption factor to re-encrypt *ES*, and re-encrypts *Voting Vector* which consists of encrypted voting content and encrypted *ES*.
- The last center recovers a voter's *ES* and compute the voting result.

ES-Center

- ES-center takes a valid voter list, and checks whether a voter is a valid voter or not through one-way untappable channel.
- He generates *ES* jointly with the last mix-center.

Bulletin Board *BB*

- Anyone can see contents of *BB*, but can not modify or erase it.

Overwritable Bulletin board *OBB*

- Only each mix-center overwrites contents in *OBB*. Other people can only see it.

2.2 Overview of e-voting

Our e-voting protocol runs as follows.

Issue of *ES*

1. We suppose that ES-center takes a valid voters list. ES-center and the last mix-center jointly generate *ES*.
2. After ES-center checks a voter's id and signature through one-way untappable channel, he sends *ES* and encrypted *ES* to a valid voter.
3. ES-center posts a valid voter's id to *BB*.

Voting stage

1. A voter chooses a voting content, and encrypts it with *ES*.
2. A voter generates *Voting Vector* which consists of encrypted voting content and encrypted *ES* by ES-center, and sends it to *OBB*.
3. The first mix-center gets *Voting Vector* from *OBB* and re-encrypts *Voting Vector* with his random encryption factor as the original universal re-encryption mix-net. He overwrites the old *Voting Vector* in *OBB* in a random order.

4. To prove a valid of mixing of the first mix-center, the first mix-center (Prover) proves to the second mix-center (Verifier) without leaking his random encryption factor (See appendix A). He sends his proof to the designated field of *BB*.

5. Other mix-centers from the second mix-center to $n-1$ mix-center re-encrypt *Voting Vector* with their random encryption factors and overwrite the old *Voting Vector* in *OBB* in order. Each mix-center proves his mixing to the next mix-center using the modified designated-verifier re-encryption proof.

Counting stage

1. The last mix-center decrypts a voter's *ES*. He computes the voting result with *ES*.
2. ES-center verifies the computed voting result with the number of issued *ES* and the published voting result by the last mix-center.

Table 1. Comparison of mix-net e-voting systems

Property	[PIK93]	[SK95]	[HS00]	Our scheme
Receipt-freeness	No	Yes	Yes	Yes
Universal Verifiability	No	Yes	Yes	Yes
Privacy	Yes	No	Yes	Yes
Robustness	Yes	No	Yes	Yes

3 Conclusion

We propose mix net e-voting system which satisfies robustness and privacy as well as receipt-freeness and universal verifiability as Table1. We introduce firstly E-voting Sheet, and use Overwritable Bulletin Board. For achieving universal verifiability

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