Inductive Inference with Bounded Mind Changes

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Abstract. In this paper, we deal with inductive inference for a class of recursive languages with a bounded number of mind changes. We introduce an n-bounded finite tell-tale and a pair of n-bounded finite tell-tales of a language, and present a necessary and sufficient condition for a class to be inferable with bounded mind changes, when the equivalence of any two languages in the class is effectively decidable. We also show that the inferability of a class from positive data strictly increases, when the allowed number of mind changes increases. In his previous paper, Mukouchi gave necessary and sufficient conditions for a class of recursive languages to be finitely identifiable, that is, to be inferable without any mind changes from positive or complete data. The results we present in this paper are natural extensions of the above results.

1. Introduction

Inductive inference is a process of hypothesizing a general rule from examples. As a correct inference criterion for inductive inference of formal languages and models of logic programming, we have mainly used Gold's identification in the limit[5]. In this criterion, an inference machine is allowed to change its guesses finitely many times, and the guesses are required to converge to a correct guess. Angluin[1], Wright[11] and Sato&Umayahara[6] discussed conditions for a class of formal languages to be inferable from positive data. Shinohara[7, 8] also discussed inductive inferability from positive data in more general setting and exhibited that inductive inference from positive data is much more powerful than it has been believed.

Considering ordinary learning process of human beings, the criterion of identification in the limit seems to be natural. However, we can not decide in general whether a sequence of guesses from an inference machine converges or not at a certain time, and the results of the inference necessarily involve some risks. In his previous paper[10], Mukouchi gave necessary and sufficient conditions for a class of recursive languages to be finitely identifiable, that is, to be inferable without any mind changes from positive or complete data. We use the phrase 'mind change' to mean that an inference machine changes its guess.

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In this paper, we deal with inductive inference for a class of recursive languages with a bounded number of mind changes. The results we present in this paper are natural extensions of the above results concerning finite identification.

Note that Case&Smith[3] discussed inductive inference of a class of recursive functions from view point of anomalies and mind changes, and showed that there is a natural hierarchy. Case&Lynes[4] also showed that an anomaly hierarchy exists even in case of a class of recursive languages.

In Section 2, we prepare some necessary concepts for our discussions. We also recall the results on finite identification from positive and complete data. In Section 3, we discuss conditions for a class to be inferable with bounded mind changes from positive data. Angluin[1] introduced the notion of a finite tell-tale of a language to discuss inferability of formal languages from positive data, and showed that a class is inferable from positive data if and only if there is a recursive procedure to enumerate all elements in the finite tell-tale of any language of the class. In this paper, we introduce an n-bounded finite tell-tale of a language, and present a necessary and sufficient condition for a class to be inferable with bounded mind changes, when the equivalence of any two languages in the class is effectively decidable. We also exhibit a concrete class of recursive languages which is inferable with at most $n$ mind changes but not inferable with at most $n - 1$ mind changes, and show that the inferability of a class strictly increases, when the allowed number of mind changes increases. Case&Smith[3] showed similar results for a class of recursive functions. In Section 4, we give a necessary and sufficient condition for a class to be inferable with bounded mind changes from complete data, which is analogous to the above condition concerning positive data.

2. Preliminaries

Let $U$ be a recursively enumerable set to which we refer as a universal set. Then we call $L \subseteq U$ a language. We do not consider the empty language in this paper.

**Definition 2.1.** A class of languages $\Gamma = L_1, L_2, \ldots$ is said to be an indexed family of recursive languages if there exists a computable function $f : N \times U \to \{0, 1\}$ such that

$$f(i, w) = \begin{cases} 1, & \text{if } w \in L_i, \\ 0, & \text{otherwise.} \end{cases}$$

From now on, we assume a class of languages is an indexed family of recursive languages without any notice.

**Definition 2.2.** A positive presentation of a language $L$ is an infinite sequence $w_1, w_2, \ldots$ of elements of $U$ such that $\{w_1, w_2, \ldots\} = L$.

A complete presentation of a language $L$ is an infinite sequence $(w_1, t_1), (w_2, t_2), \ldots$ of elements of $U \times \{0, 1\}$ such that $\{w_i \mid t_i = 1, i \geq 1\} = L$ and $\{w_j \mid t_j = 0, j \geq 1\} = U - L$.

Positive or complete presentations are denoted by $\sigma, \delta$, the finite sequence which consists of first $n \geq 0$ data in $\sigma$ by $\sigma[n]$ and the finite set by $\sigma(n)$.

For a finite sequence $\sigma[n]$ and a sequence $\delta$, the sequence which is obtained by concatenating $\sigma[n]$ with $\delta$ is denoted by $\sigma[n] \cdot \delta$. 

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**Definition 2.3.** An $n$-bounded inference machine (abbreviated to $IM_n$; $n \geq 0$ or $n = \ast$) is an effective procedure that requests inputs from time to time and produces outputs from time to time, where if $n \geq 0$, it produces at most $n + 1$ outputs, and if $n = \ast$, it produces at most finitely many outputs.

The outputs produced by the machine are called guesses.

For a finite sequence $\sigma[m] = w_1, w_2, \ldots, w_m$, we denote by $M(\sigma[m])$ the last guess produced by an $IM_n$ $M$ which is successively fed $w_1, w_2, \ldots, w_m$ on its input requests.

The inference machines we are dealing with in this paper may not produce a guess after reading a datum until requesting a next datum.

**Definition 2.4.** A class $\Gamma = L_1, L_2, \cdots$ is said to be $EX_n$ identifiable from positive data (resp., complete data) if there exists an $IM_n$ $M$ satisfying the following ($n \geq 0$ or $n = \ast$): For any language $L_i$ of $\Gamma$ and for any positive presentation (resp., complete presentation) $\sigma$ of $L_i$, the last guess $k$ of $M$ which is successively fed $\sigma$'s data satisfies $L_k = L_i$.

A class $\Gamma$ is said to be finitely identifiable (resp., identifiable in the limit) if it is $EX_0$ identifiable (resp., $EX_\ast$ identifiable).

A class $\Gamma$ is also said to be $EX$-TXT$_n$ identifiable (resp., $EX$-INF$_n$ identifiable) if it is $EX_n$ identifiable from positive data (resp., complete data). By the same notation $EX$-TXT$_n$ (resp., $EX$-INF$_n$), we also denote the set of the classes that are $EX$-TXT$_n$ identifiable (resp., $EX$-INF$_n$ identifiable).

In this paper, a finite-set-valued function $F$ is said to be computable if there exists an effective procedure that produces all elements in $F(x)$ and then halts uniformly for any argument $x$.

Mukouchi[10] presented necessary and sufficient conditions for an indexed family of recursive languages to be finitely identifiable.

**Definition 2.5 (Mukouchi[10]).** A set $S_i$ is said to be a definite finite tell-tale of $L_i$ if

1. $S_i$ is a finite subset of $L_i$, and
2. $S_i \subseteq L_j$ implies $L_j = L_i$ for any index $j$.

**Theorem 2.1 (Mukouchi[10]).** A class $\Gamma$ is finitely identifiable from positive data if and only if a definite finite tell-tale of $L_i$ is uniformly computable for any index $i$, that is, there exists an effective procedure that on input $i$ produces all elements of a definite finite tell-tale of $L_i$ and then halts.

**Definition 2.6 (Mukouchi[10]).** A language $L$ is said to be consistent with a pair of sets $(T_i, F_i)$ if $T_i \subseteq L$ and $F_i \subseteq U - L$.

A pair of sets $(T_i, F_i)$ is said to be a pair of definite finite tell-tales of $L_i$ if

1. $T_i$ is a finite subset of $L_i$,
2. $F_i$ is a finite subset of $U - L_i$, and
3. if $L_j$ is consistent with the pair $(T_i, F_i)$, then $L_j = L_i$.

**Theorem 2.2 (Mukouchi[10]).** A class $\Gamma$ is finitely identifiable from complete data if and only if a pair of definite finite tell-tales of $L_i$ is uniformly computable for any index $i$.
The following corollary shows a necessary condition for a class to be finitely identifiable.

**Corollary 2.3.** If a class \( I' \) is finitely identifiable from positive or complete data, then whether \( L_i = L_j \) or not is effectively decidable for any indices \( i, j \).

**Proof:** Clearly, if \( I' \) is finitely identifiable from positive data, then \( I' \) is also finitely identifiable from complete data. Therefore, it suffices to show the case of complete data.

Suppose \( I' \) is finitely identifiable from complete data. Fix arbitrary indices \( i, j \). To begin with, compute a pair of definite finite tell-tales of \( L_i \), and set it to \( (T_i, F_i) \). We can effectively compute this pair by Theorem 2.2. Then check whether \( L_j \) is consistent with \( (T_i, F_i) \). We can effectively check this, because \( T_i \) and \( F_i \) are explicitly given finite sets. If \( L_j \) is not consistent with \( (T_i, F_i) \), then we conclude \( L_i \neq L_j \), because \( L_i \) is consistent with \( (T_i, F_i) \). Otherwise, we conclude \( L_i = L_j \) by Definition 2.6. ■

### 3. Inductive Inference with Bounded Mind Changes from Positive Data

First of all, we give a necessary condition for a class \( I' \) to be \( EX\text{-}TXT_n \) identifiable.

**Proposition 3.1.** For any \( n \geq 1 \), if a class \( I' \) contains languages \( L_{i_0}, L_{i_1}, \ldots, L_{i_n} \) such that \( L_{i_0} \subseteq L_{i_1} \subseteq \cdots \subseteq L_{i_n} \), then \( I' \) is not \( EX\text{-}TXT_{n-1} \) identifiable.

**Proof:** Suppose that \( I' \) contains languages \( L_{i_0}, L_{i_1}, \ldots, L_{i_n} \) such that \( L_{i_0} \subseteq L_{i_1} \subseteq \cdots \subseteq L_{i_n} \) and that \( I' \) is \( EX\text{-}TXT_1 \) identifiable by an IM, \( M \). For simplicity, put \( L_j = L_{i_j} \) (\( 0 \leq j \leq n \)). We show that \( M \) needs to change its guesses more than \( n \) times to identify \( L'_n \) from a certain positive presentation of \( L'_n \). Let \( \sigma_j \) be an arbitrary positive presentation of \( L'_j \). We recursively define \( c_j \) and \( \delta_j \) as follows:

**Stage 0:**
- Let \( c_0 := 0 \) and \( \delta_0 := \sigma_0 \). Goto Stage 1.

**Stage \( m(1 \leq m \leq n) \):**
- Let \( c_m := \min \{ c > c_{m-1} \mid \exists k \text{ s.t. } M(\delta_{m-1}[c]) = k \wedge L'_{m-1} = L_k \} \) and \( \delta_m := \delta_{m-1}[c_m] \cdot \sigma_m \).
  - Such an integer \( c_m \) exists, because \( \delta_{m-1} \) is a positive presentation of \( L'_{m-1} \) and \( I' \) is \( EX\text{-}TXT_1 \) identifiable by \( M \). Note that the above \( \delta_m \) becomes a positive presentation of \( L'_m \), because \( L'_{m-1} \subseteq L'_m \).
  - Goto Stage \( m + 1 \).

**Stage \( n + 1 \):**
- Let \( c_{n+1} := \min \{ c > c_n \mid \exists k \text{ s.t. } M(\delta_n[c]) = k \wedge L'_n = L_k \} \).

When we feed a positive presentation \( \delta_n \) of \( L'_n \) successively to \( M \), it should output guesses after reading \( c_1 \)-th, \( c_2 \)-th, \ldots, \( c_{n+1} \)-th datum, and so it can not identify \( L'_n \) within \( n \) mind changes. ■

Before proceeding to the next corollary, we briefly recall a pattern and a pattern language. (For more details, see Angluin[2] or Mukouchi[9].)
Fix a finite alphabet with at least two constant symbols. A pattern is a nonnull finite string of constant and variable symbols. The pattern language \( L(\pi) \) generated by a pattern \( \pi \) is the set of all strings obtained by substituting nonnull strings of constant symbols for the variables of \( \pi \). Since two patterns that are identical except for renaming of variables generate the same pattern language, we do not distinguish one from the other. We can enumerate all patterns recursively and whether \( w \in L(\pi) \) or not for any \( w \) and \( \pi \) is effectively decidable. Therefore, we can consider the class of pattern languages as an indexed family of recursive languages, where the pattern itself is considered to be an index.

**Corollary 3.2.** For any \( n \geq 0 \), the class of pattern languages is not \( \text{EX-TXT}_n \) identifiable.

**Proof:** By Proposition 3.1, it suffices to show that there exist patterns \( \pi_0, \pi_1, \ldots, \pi_m \) such that \( L(\pi_0) \subset L(\pi_1) \subset \cdots \subset L(\pi_m) \) for any \( m \geq 1 \).

In fact, let \( \pi_0 = x_1 x_2 \cdots x_{m+1}, \pi_1 = x_1 x_2 \cdots x_m, \ldots, \pi_m = x_1 \). Then \( L(\pi_i) \) is the set of all constant strings of length more than \( m - i \), and clearly
\[
L(\pi_0) \subset L(\pi_1) \subset \cdots \subset L(\pi_m).
\]

Angluin[1] showed that the class of pattern languages is inferable from positive data in the limit, that is, it is \( \text{EX-TXT}_n \) identifiable.

**Definition 3.1.** Let \( \Gamma = L_1, L_2, \ldots \). A set \( S_i \) is said to be a 0-bounded finite tell-tale (abbreviated to \( \text{FT}_0 \)) in \( \Gamma \) of \( L_i \) if \( S_i \) is a definite finite tell-tale of \( L_i \), that is,

1. \( S_i \) is a finite subset of \( L_i \), and
2. \( S_i \subseteq L_j \) implies \( L_j = L_i \) for any index \( j \).

A set \( S_i \) is said to be an \( n \)-bounded finite tell-tale (abbreviated to \( \text{FT}_n ; n \geq 1 \)) in \( \Gamma \) of \( L_i \) if

1. \( S_i \) is a finite subset of \( L_i \), and
2. if \( L_j \neq L_i \) and \( S_i \subseteq L_j \), then there exist an \( \text{FT}_{n-1} \) in \( \Gamma \) of \( L_j \).

Intuitively, an \( \text{FT}_n \) in \( \Gamma \) of \( L_i \) is a tell-tale which validates producing the guess \( i \), when the inference machine is allowed to produce another \( n - 1 \) guesses.

We can easily prove by induction on \( n \) that if a certain finite set \( S \) is an \( \text{FT}_n \) in \( \Gamma \) of \( L_i \), then \( S \) is also an \( \text{FT}_{n+1} \) in \( \Gamma \) of \( L_i \) (\( n \geq 0 \)).

**Definition 3.2.** An \( \text{FT}_0 \) in \( \Gamma \) of \( L_i \) is said to be recurrently computable if a certain \( \text{FT}_0 S_i \) in \( \Gamma \) of \( L_i \) is computable.

An \( \text{FT}_n \) in \( \Gamma \) of \( L_i \) is said to be recurrently computable (\( n \geq 1 \)) if

1. a certain \( \text{FT}_n S_i \) in \( \Gamma \) of \( L_i \) is computable, and
2. for any index \( j \), if \( L_j \neq L_i \) and \( S_i \subseteq L_j \), then an \( \text{FT}_{n-1} \) in \( \Gamma \) of \( L_j \) is recurrently computable.

An \( \text{FT}_n \) of \( \Gamma \) is said to be recurrently constructible if an \( \text{FT}_n \) in \( \Gamma \) of \( L_i \) is recurrently computable for any index \( i \) (\( n \geq 0 \)).
Lemma 3.3. Suppose whether \( L_i = L_j \) or not is effectively decidable for any indices \( i,j \).

For any \( n \geq 0 \), if a class \( \Gamma \) is \( EX-TXT_n \) identifiable, then an \( FT_n \) of \( \Gamma \) is recurrently constructible.

Proof: Suppose \( \Gamma \) is \( EX-TXT_n \) identifiable by an \( IM_n \) \( M \). In what follows, for a finite sequence \( \varphi \), we denote the corresponding finite set by \( \bar{\varphi} \).

We consider the following partial recursive procedure which produces a finite sequence of \( U \): 

Procedure \( Ft(m, i) \);
begin
    for each finite sequence \( \psi \) of \( U \) do
        /* Note that all finite sequences of \( U \) are recursively enumerable */
        if \( \bar{\psi} \subseteq L_i \) do begin
            Initialize \( M \);
            Feed successively \( \psi \) to \( M \) on its input requests;
            if \( M \) produces any guess then begin
                Let \( k \) be the number of guesses \( (1 \leq k \leq n + 1) \) and \( g \) be the last guess produced by \( M \);
                if \( (k > n - m) \land (L_g = L_i) \) then output \( \psi \) and stop;
            end;
        end;
    end.

(1) Clearly, if \( Ft(m, i) \) is defined, then \( \bar{Ft}(m, i) \) is a finite subset of \( L_i \).

(2) If \( Ft(0, i) \) is defined, then \( \bar{Ft}(0, i) \) is an \( FT_0 \) of \( L_i \). In fact, suppose that \( Ft(0, i) \) is defined and that \( \bar{Ft}(0, i) \) is not an \( FT_0 \) of \( L_i \). Then, there exists an index \( j \) such that \( L_j \neq L_i \) and \( \bar{Ft}(0, i) \subseteq L_j \). Let \( \sigma_j \) be an arbitrary positive presentation of \( L_j \). Since \( \bar{Ft}(0, i) \cdot \sigma_j \) is a positive presentation of \( L_j \), it follows that \( M \) can not identify \( L_j \) from \( Ft(0, i) \cdot \sigma_j \) with at most \( n \) mind changes, which contradicts the assumption.

(3) For any \( m \) \((1 < m \leq n)\) and any index \( i \), if \( Ft(m, i) \) is defined and there exists an index \( j \) such that \( L_j \neq L_i \) and \( \bar{Ft}(m, i) \subseteq L_j \), then \( Ft(m - 1, j) \) is defined. In fact, suppose \( Ft(m - 1, j) \) is not defined. Let \( \sigma_j \) be an arbitrary positive presentation of \( L_j \) and \( \delta = Ft(m, i) \cdot \sigma_j \). Then there exists a \( k > |Ft(m, i)| \) such that \( M(\delta[k]) = g \) and \( L_g = L_j \) for some \( g \). Since \( \delta[k] \) is a finite sequence of \( U \), it should appear in the for loop above. Furthermore, when \( M \) is successively fed \( \delta[k] \) on its input requests, it should produce more than \( n - m + 1 \) guesses. Hence the procedure will produce an output. This is a contradiction.

By (1), (2) and (3), for any \( m \) \((0 < m \leq n)\) and any index \( i \), if \( Ft(m, i) \) is defined, then it is an \( FT_m \) in \( \Gamma \) of \( L_i \).

Moreover, for any index \( i \), \( Ft(n, i) \) is defined, since \( \Gamma \) is identifiable by \( M \). Therefore, an \( FT_n \) of \( \Gamma \) is recurrently constructible.
Lemma 3.4. Suppose whether \( L_i = L_j \) or not is effectively decidable for any indices \( i, j \).

For any \( n \geq 0 \), if an \( FT_n \) of a class \( \Gamma \) is recurrently constructible, then \( \Gamma \) is EX-TXT\(_n\) identifiable.

Proof: Suppose an \( FT_n \) of \( \Gamma \) is recurrently constructible. We denote by \( FT_m(i) \) the result of computation of an \( FT_m \) in \( \Gamma \) of \( L_i \) \((m \geq 0)\). We consider the following procedure:

Procedure \( M \);
begin
\( m := n; \quad k := 0; \)
\( S := \phi; \quad T := \phi; \)
for \( j := 1 \) to \( \infty \) do begin
   read a next datum and add it to \( T \); /* Note that \( \sharp T = j */
   for \( i := 1 \) to \( j \) do
     if \((k = 0) \lor (L_k \neq L_i \land S \subseteq L_i)\) then
       if \( FT_m(i) \subseteq T \) then begin
         output \( i \);
         if \( m = 0 \) then stop;
         \( S := FT_m(i); \quad k := i; \)
         \( m := m - 1; \)
       end;
   end;
end;

Clearly, this procedure produces at most \( n + 1 \) outputs. Suppose we are going to feed a positive presentation \( \sigma \) successively to the procedure on its input requests.

(1) This procedure produces at least one guess. In fact, suppose this procedure never produces a guess. When it reaches the case

\[ j = \max\{h, \min\{l \mid FT_n(h) \subseteq \sigma(l)\}\} \quad \text{and} \quad i = h, \]

this procedure should produce the guess \( h \), which contradicts the assumption.

(2) Suppose the last guess, say \( g \), produced by this procedure is not correct.

(i) In case of \( m = 0 \), when the procedure produced the last guess. It contradicts the definition of an \( FT_0 \).

(ii) Otherwise, note that \( S \subseteq T \) and \( L_g \neq L_h \). When it reaches the case

\[ j \geq \max\{h, \min\{l \mid FT_n(h) \subseteq \sigma(l)\}\} \quad \text{and} \quad i = h, \]

this procedure should produce the next guess \( h \), which contradicts the assumption. \( \square \)

We obtain the following Theorem 3.5 by Lemma 3.3 and Lemma 3.4.

Theorem 3.5. Suppose whether \( L_i = L_j \) or not is effectively decidable for any indices \( i, j \).

For any \( n \geq 0 \), a class \( \Gamma \) is EX-TXT\(_n\) identifiable if and only if an \( FT_n \) of \( \Gamma \) is recurrently constructible.
Note that in case of $n = 0$, the above theorem is equivalent to Theorem 2.1 by Corollary 2.3.

**Example 3.1.** Fix an arbitrary number $n \geq 0$. We consider the following set:

$$C_n = \left\{ (q_1, q_2, \ldots, q_{n+1}) \in N^{n+1} \mid q_1, q_2, \ldots, q_{n+1} \text{ are prime numbers with } q_1 < q_2 < \cdots < q_{n+1} \right\}.$$

Fix an arbitrary computable bijection from $C_n$ to $N$, and we denote it by $\langle \rangle$. We consider the following class:

$$\Gamma = L_1, L_2, \ldots,$$

where

$$L_{\langle q_1, q_2, \ldots, q_{n+1} \rangle} = \{ m \in N \mid m \text{ is a multiple of } q_j \text{ for some } j \ (1 \leq j \leq n+1) \}.$$  

Clearly, this class $\Gamma$ is an indexed family of recursive languages. This class is also finitely identifiable. In fact, there exists a computable $FT_0$ of $L_{\langle q_1, q_2, \ldots, q_{n+1} \rangle}$ such as

$$\{q_1, q_2, \ldots, q_{n+1}\}.$$

Note that if $n \geq 1$, this class $\Gamma$ does not have the property of so-called finite thickness [11, 11], which is a sufficient condition to be inferable from positive data in the limit.

**Example 3.2.** Fix an arbitrary number $n \geq 0$. We consider the following set:

$$D_n = \left\{ (q_1, q_2, \ldots, q_k) \in N^k \mid 1 \leq k \leq n+1, q_1, q_2, \ldots, q_k \text{ are prime numbers with } q_1 < q_2 < \cdots < q_k \right\}.$$  

Fix an arbitrary computable bijection from $D_n$ to $N$, and we denote it by $\llbracket \rrbracket$. We consider the following class:

$$\Gamma' = L_1, L_2, \ldots,$$

where

$$L_{\llbracket q_1, q_2, \ldots, q_k \rrbracket} = \{ m \in N \mid m \text{ is a multiple of } q_j \text{ for some } j \ (1 \leq j \leq k) \} \quad (1 \leq k \leq n+1).$$

Clearly, this class $\Gamma'$ is an indexed family of recursive languages. This class is also $EX$-$TXT_n$ identifiable. In fact, there exists a computable $FT_m$ ($n - k + 1 \leq m \leq n$) of $L_{\llbracket q_1, q_2, \ldots, q_k \rrbracket}$ such as

$$FT_{n-k+1} = \cdots = FT_n = \{q_1, q_2, \ldots, q_k\},$$

and it follows that an $FT_n$ of $\Gamma'$ is recurrently constructible.

On the other hand, this class is shown to be not $EX$-$TXT_{n-1}$ identifiable by Proposition 3.1 if $n \geq 1$.

Note that this class $\Gamma'$ does not have the property of finite thickness if $n \geq 1$.

From Corollary 3.2, Example 3.2 and the fact that the class of pattern languages is $EX$-$TXT_0$ identifiable but not $EX$-$TXT_n$ identifiable for any $n \geq 0$, we see that there exists a hierarchy such as

$$EX$-$TXT_0 \subset EX$-$TXT_1 \subset \cdots \subset EX$-$TXT_n \subset \cdots \subset EX$-$TXT_\omega.$$
4. Inductive Inference with Bounded Mind Changes from Complete Data

In this section, we give the results concerning complete data, which are analogous to the results in Section 3 concerning positive data.

The following Definition 4.1 and Theorem 4.1 form a remarkable contrast to Definition 3.1 and Theorem 3.5 concerning positive data.

**Definition 4.1.** Let \( \Gamma = L_1, L_2, \ldots \). A pair \( \langle T_i, F_i \rangle \) is said to be a pair of 0-bounded finite tell-tales (abbreviated to \( \text{PFT}_0 \)) in \( \Gamma \) of \( L_i \) if \( \langle T_i, F_i \rangle \) is a pair of definite finite tell-tales of \( L_i \), that is,

1. \( T_i \) is a finite subset of \( L_i \),
2. \( F_i \) is a finite subset of \( U - L_i \), and
3. for any index \( j \), if \( L_j \) is consistent with the pair \( \langle T_i, F_i \rangle \), then \( L_j = L_i \).

A pair \( \langle T_i, F_i \rangle \) is said to be a pair of \( n \)-bounded finite tell-tales (abbreviated to \( \text{PFT}_n \); \( n \geq 1 \)) in \( \Gamma \) of \( L_i \) if

1. \( T_i \) is a finite subset of \( L_i \),
2. \( F_i \) is a finite subset of \( U - L_i \), and
3. for any index \( j \), if \( L_j \neq L_i \) and \( L_j \) is consistent with \( \langle T_i, F_i \rangle \), then there exists a \( \text{PFT}_{n-1} \) in \( \Gamma \) of \( L_j \).

We can easily prove by induction on \( n \) that if a certain pair of finite sets \( (T, F) \) is a \( \text{PFT}_n \) in \( \Gamma \) of \( L_i \), then \( \langle T, F \rangle \) is also a \( \text{PFT}_{n+1} \) in \( \Gamma \) of \( L_i \) (\( n \geq 0 \)).

Similarly to Definition 3.2, we define the recurrent computability and the recurrent constructibility of a \( \text{PFT}_n \).

We can also prove the following theorem in a similar way to the proofs of Lemma 3.3 and Lemma 3.4.

**Theorem 4.1.** Suppose whether \( L_i = L_j \) or not is effectively decidable for any indices \( i, j \).

For any \( n \geq 0 \), a class \( \Gamma \) is \( \text{EX-INF}_n \) identifiable if and only if a \( \text{PFT}_n \) of \( \Gamma \) is recurrently constructible.

Note that in case of \( n = 0 \), the above theorem is equivalent to Theorem 2.2 by Corollary 2.3.

5. Concluding Remarks

We have investigated some characterization theorems on language learning with a bounded number of mind changes, and exhibited necessities of mind changes.

We have also recognized again the importance of paying attention to a characteristic subset of each language in the class, when we consider language learning.
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