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Shoudai, Takayoshi

Department of Control Engineering and Science Kyushu Institute of Technology

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A P-Complete Language Describable with Iterated Shuffle

Takayoshi Shoudai
Department of Control Engineering and Science
Kyushu Institute of Technology
Iizuka 820, Japan

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Abstract

We show that a P-complete language can be described by using the shuffle operator, shuffle closure, union, concatenation, Kleene star and intersection on a finite alphabet.

1 Introduction

In this paper, we construct a P-complete language by using shuffle operator Δ , iterated shuffle \dagger , union \cup , concatenation \cdot , Kleene star $*$ and intersection \cap over a finite alphabet. The shuffle operator was introduced by [10] to describe the class of flow expressions. Formal properties of expressions with these operators have been extensively studied from various points in the literatures [2, 3, 4, 5, 8, 9, 10, 11].

It is known that the complexity of almost classes of languages can be increased by using the iterated shuffle operator. For example, there are two deterministic context-free languages L_1 and L_2 such that $L_1 \Delta L_2$ is NP-complete [9]. Moreover, by allowing the synchronization mechanisms, any recursively enumerable set can be described [1, 3].

In [2, 11], by using the shuffle and iterated shuffle operators together with $\cup, \cdot, *, \cup$, an NP-complete language is described. We employ the same set of operators to describe our P-complete language. In the proof of P-completeness, the intersection operator plays an important role to make the language polynomial-time recognizable. However, we do not know whether the intersection operator is necessary to define a P-complete language as in the case with NP-complete [2, 11].

Recently, P-complete problems have received considerable attentions since they do not seem to allow any efficient parallel algorithms [7]. This paper gives a P-complete problem of a new kind, which is described by a single expression.

2 Preliminaries

Let Σ be a finite alphabet and Σ^* be $\{a_1 \cdots a_n \mid a_i \in \Sigma \text{ for } i = 1, \dots, n \text{ and } n \geq 0\}$. A subset of Σ^* is called a *language*.

Definition 1 For languages L , L_1 and L_2 , we define the *shuffle operator* Δ , the *iterated shuffle* \dagger and operators, $\cdot, *, +$ as follows:

- (1) $L_1 \Delta L_2 = \{x_1 y_1 x_2 y_2 \cdots x_m y_m \mid x = x_1 x_2 \cdots x_m \in L_1, y = y_1 y_2 \cdots y_m \in L_2 \text{ and } x_i, y_i \in \Sigma^* \text{ for } i = 1, \dots, m\}$ (shuffle operator).
- (2) $L^\dagger = \{\varepsilon\} \cup L \cup (L \Delta L) \cup (L \Delta L \Delta L) \cup \cdots$ (iterated shuffle).
- (3) $L_1 \cdot L_2 = \{xy \mid x \in L_1 \text{ and } y \in L_2\}$ (abbreviated to $L_1 L_2$).
- (4) $L^* = \{\varepsilon\} \cup L \cup (L \cdot L) \cup (L \cdot L \cdot L) \cdots$.
- (5) $L^+ = L \cdot L^*$.

We identify a language $\{w\}$ which consists of only one word with the w . Thus, we will denote $\{w\}^*, \{w\}^+, \{w\}^\dagger, \dots$ by w^*, w^+, w^\dagger , respectively.

As the basis of our reduction, we use the circuit value problem (CVP) that was shown P-complete [6]. Our definition in this paper slightly different from one in [6].

CIRCUIT VALUE PROBLEM (CVP)

INSTANCE: A circuit $C = (C_1, \dots, C_m, C_{m+1}, \dots, C_n)$, where each C_i is either (i) $C_i = \text{true}$ or *false* ($1 \leq i \leq m$), (ii) $C_i = \text{NOR}(C_j, C_k)$ ($m+1 \leq i \leq n$ and $j, k < i$).

PROBLEM: Decide whether the value of C_n is *true*.

In later section, CVP represents the set of all circuits whose output is *true*.

Let Σ be a finite alphabet, v_1, v_2, \dots, v_m be symbols where $v_i \in \Sigma$ for $i = 1, \dots, m$ and w_1, w_2, \dots, w_{m+1} be words on a alphabet $\Sigma - \{v_1, v_2, \dots, v_m\}$. By using the iterated shuffle operation, a language $\{v_1^n v_2^n \cdots v_m^n \mid n \geq 1\}$ can be described as $(v_1 v_2 \cdots v_m)^\dagger \cap v_1^+ v_2^+ \cdots v_m^+$. Moreover, we can represent $\{w_1 v_1^n w_2 v_2^n \cdots w_m v_m^n w_{m+1} \mid n \geq 1\}$ as

$$(w_1 w_2 \cdots w_{m+1} \Delta (v_1 v_2 \cdots v_m)^\dagger) \cap w_1 v_1^+ w_2 v_2^+ \cdots w_m v_m^+ w_{m+1}.$$

We often use this form of languages to define a P-complete language. Whenever languages like these are defined in the next section, we will not describe the languages explicitly by using the shuffle operation and the iterated shuffle.

3 A P-complete language

The main result in this paper is the following theorem.

Theorem 1 A P-complete language can be described with operators $\cdot, *, \cup, \cap, \Delta, \dagger$.

3.1 Definition of the language

We will describe a P-complete language \mathcal{L} with the alphabet $\Sigma = \{0, 1, a, b, c, d, u, v, x, y\}$. This language is defined stepwize.

At first, a language L is defined as follows:

$$\begin{aligned}
 L_a &= a^+0 \cup a^+1 = \{a^i\beta \mid i \geq 1 \text{ and } \beta \in \{0, 1\}\}. \\
 L_{bba} &= (b^+1b^+1a^+0) \cup (b^+0b^+1a^+1) \cup (b^+1b^+0a^+1) \cup (b^+0b^+0a^+1) \\
 &= \{b^j\beta'b^k\beta''a^i\beta \mid i, j, k \geq 1 \text{ and } (\beta', \beta'', \beta) \in \{(1, 1, 0), (0, 1, 1), (1, 0, 1), (0, 0, 1)\}\} \\
 L_b &= b^+1 = \{b^i1 \mid i \geq 1\}. \\
 \hline
 L &= cL_a^+L_{bba}^+L_b.
 \end{aligned}$$

The following language T (resp. F) is used for a distribution of *true* (resp. *false*) value.

$$\begin{aligned}
 T_x &= \{1dx^iu^i \mid i \geq 1\}, \quad T_y = \{1y^iv^i \mid i \geq 1\}. \\
 T_{xy} &= \{1dx^iu^i1y^iv^i \mid i \geq 1\}, \quad T_{yy} = \{1y^iv^i1y^iv^i \mid i \geq 1\}. \\
 \hline
 T_{odd} &= T_{xy}T_{yy}^*T_y \cap T_xT_{yy}^* = \{1dx^iu^i(1y^iv^i)^j \mid i \geq 1, j \geq 1 \text{ and } j \text{ is odd}\}. \\
 T_{even} &= T_xT_{yy}^*T_y \cap T_{xy}T_{yy}^* = \{1dx^iu^i(1y^iv^i)^j \mid i \geq 1, j \geq 1 \text{ and } j \text{ is even}\}. \\
 \hline
 T &= T_x \cup T_{odd} \cup T_{even} = \{1dx^iu^i(1y^iv^i)^j \mid i \geq 1 \text{ and } j \geq 0\}.
 \end{aligned}$$

F is defined in a similar way. We use a symbol 0 instead of 1 which is used to construct the language T .

$$F = \{0dx^iu^i(0y^iv^i)^j \mid i \geq 1 \text{ and } j \geq 0\}.$$

Subwords $1y^iv^i$ (resp. $0y^iv^i$) of a word in T (resp. F) are combined with b^i0 (resp. b^i1) of words in L and decides the value of the i th variable. These three languages L , T and F are combined with each other by using the shuffle operation and the iterated shuffle.

$$\mathcal{J} = L\Delta(T \cup F)^\dagger.$$

A language \mathcal{K} is used for our language to become polynomial time decidable. We construct the language \mathcal{K} stepwise as follows:

$$\begin{aligned} A_{11} &= \{a^i 11 dx^i u^i \mid i \geq 1\}. \\ A_{00} &= \{a^i 00 dx^i u^i \mid i \geq 1\}. \\ A_{01} &= \{a^i 01 dx^i u^i \mid i \geq 1\}. \end{aligned}$$

In a similar way, following languages are defined.

$$\begin{aligned} B_{01} &= \{b^i 01 y^i v^i \mid i \geq 1\}. \\ B_{11} &= \{b^i 11 y^i v^i \mid i \geq 1\}. \end{aligned}$$

$$\overline{M = (A_{11} \cup A_{00})^+ (B_{01} B_{01} A_{01})^+ B_{11}}.$$

The language M has words whose subwords of the form $dx^i u^i$ corresponding to the i th gate occurred more than two times a word. We want these subwords to be occurred exactly one time a word.

$$\begin{aligned} N_d &= (dxudx^2u^2\Delta(xuru)^\dagger) \cap (dx^+u^+dx^+u^+) = \{dx^i u^i dx^{i+1} u^{n+1} \mid i \geq 1\}. \\ \overline{N} &= c((dxuN_d^* \cap N_d^* dx^+ u^+) \cup (dxuN_d^* dx^+ u^+ \cap N_d^*)) \\ &= \{cdxudx^2u^2 \dots dx^i u^i \mid i \geq 1\}. \end{aligned}$$

Then, we define a language \mathcal{K} which will be used for allowing a language \mathcal{J} to be in P.

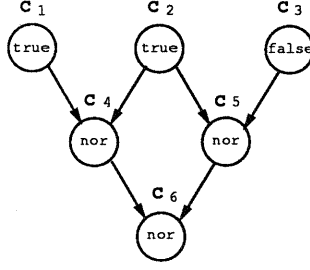
$$\mathcal{K} = M \cap (N\Delta\Sigma'), \text{ where } \Sigma' = \Sigma - \{d, u, x\}.$$

Finally, we defined a language \mathcal{L} as follows:

$$\mathcal{L} = \mathcal{J} \cap \mathcal{K}.$$

3.2 Proof of the P-completeness

Theorem 1 follows from a next lemma.



$$w = a11dxua^211dx^2u^2a^300dx^3u^3b01yvb^201y^2v^2a^401dx^4u^4 \\ b^201y^2v^2b^301y^3v^3a^501dx^5u^5b^401y^4v^4b^501y^5v^5a^601dx^6u^6b^611y^6v^6.$$

Figure 1: This circuit is transformed to a word w .

Lemma 1 \mathcal{L} is log-space equivalent to CVP, i.e., \mathcal{L} is log-space reducible from CVP and CVP is log-space reducible from \mathcal{L} .

Proof. We will define a function f from CVP to Σ^* . f is a function which transform $C = (C_1, \dots, C_n) \in \text{CVP}$ to $f(C) = \gamma w_1 \cdots w_n w_{n+1} \in \Sigma^*$, where

$$w_i = \begin{cases} a^i 11dx^i u^i & (C_i = \text{true}) \\ a^i 00dx^i u^i & (C_i = \text{false}) \\ b^j 01y^j v^j b^k 01y^k v^k a^i 01dx^i u^i & (C_i = \text{NOR}(C_j, C_k)) \\ b^n 11y^n v^n & (i = n + 1). \end{cases}$$

It is easy to see that this function is computable in log-space by using a deterministic Turing machine.

We show following two claims.

Claim 1. $f(C) \in \mathcal{L}$, for every $C \in \text{CVP}$.

Proof. Let a word $w = cw_1 \cdots w_m w_{m+1} \cdots w_n w_{n+1}$ be a transformed word from some n -gates instance $C = (C_1, \dots, C_m, C_{m+1}, \dots, C_n)$ where C_i is an *input* gate for $1 \leq i \leq m$, an *and* gate for $m + 1 \leq i \leq n$ and an output of this circuit is *true*. This instance has only one tuple of assignments of a boolean value (*true* or *false*) to each variables. We describe this assignment as $B = (\beta_1, \dots, \beta_n)$ such that $\beta_i = 1$ (resp. $\beta_i = 0$) if $C_i = \text{true}$ (resp. $C_i = \text{false}$) for $i = 1, \dots, n$.

According to $B = (\beta_1, \dots, \beta_n)$, we divide w_i into two words w_i' and w_i'' .

(1) For $i = 1, \dots, m$, $w_i' = a^i \beta_i$, $w_i'' = \beta_i dx^i u^i$.

(2) For $i = m + 1, \dots, n$, $w_i' = b^j \bar{\beta}_j b^k \bar{\beta}_k a^i \bar{\beta}_i$, $w_i'' = \beta_j y^j v^j \beta_k y^k v^k \beta_i dx^i u^i$.

We note that since C_j, C_k and C_i are related with each other by an NOR gate, w_i' is in L_{bba} .

$$(3) w_{n+1}' = b^n 1, w_{n+1}'' = 1y^n v^n.$$

It is easy to see that a word $w' = cw_1' \cdots w_{n+1}'$ is in $L = L_a^+ L_{bba}^+ L_b$.

On the other hand, since $w'' = w_1'' \cdots w_{n+1}''$ is constructed with subwords of the form $\beta_i dx^i u^i$ or $\beta_i y^i v^i$ and for each NOR gate, input gate numbers of this gate are always lower than a number of itself, we can describe the word $w'' \in t_1 \Delta t_2 \Delta \cdots \Delta t_n$, where $t_i = \beta_i dx^i u^i \beta_i y^i v^i \cdots \beta_i y^i v^i$. Since $t_i \in T$ or F , for $i = 1, \dots, n$, $f(C) = cw_1 \cdots w_m w_{m+1} \cdots w_n w_{n+1} \in w' \Delta t_1 \Delta \cdots \Delta t_n \subset L \Delta (T \cup F)^\dagger = \mathcal{L}$.
□

Since every words w of \mathcal{L} is contained in M , c is of the form $w = cw_1 \cdots w_m w_{m+1} \cdots w_n w_{n+1}$, where, for $i = 1, \dots, n + 1$,

$$w_i = \begin{cases} a^{\beta_i} \beta_i \beta_i dx^{\beta_i} u^{\beta_i} & (1 \leq i \leq m, \beta_i \in \{0, 1\}) \\ b^{\beta_i'} 0 1 y^{\beta_i'} v^{\beta_i'} b^{\beta_i''} 0 1 y^{\beta_i''} v^{\beta_i''} a^{\beta_i} 0 1 dx^{\beta_i} u^{\beta_i} & (m + 1 \leq i \leq n) \\ b^{\beta_i} 1 1 y^{\beta_i} v^{\beta_i} & (i = n + 1) \end{cases}$$

We transform a word $w \in \mathcal{L}$ to a circuit $C = (C_1, \dots, C_m, C_{m+1}, \dots, C_n)$ as follows:

- (1) For $i = 1, \dots, m$, if $\beta_i = 1$ then $C_i = true$ else $C_i = false$.
- (2) For $i = m + 1, \dots, n$, $C_i = \text{NOR}(C_j, C_k)$ where $j = \beta_i'$ and $k = \beta_i''$.

It is easy to see that g is well-defined function and this function is log-space computable.

Claim 2. $g(w) \in CVP$, for every $w \in \mathcal{L}$.

Proof. Since $w \in N$, $\ell_i = i$ for every $i = 1, \dots, n$. Moreover, since some parts of w are constructed of words which are contained in T or F , a subword $y^i v^i$ of w is never occurred before a subword $dx^i u^i$ of w . Therefore $j, k \leq i$.

Since $w \in L \Delta (T \cup F)^\dagger$ and w includes n subwords $dxu, dx^2u^2, \dots, dx^nu^n$, there exist n words t_1, \dots, t_n in $T \cup F$ which contribute a construction of w by using the iterated shuffle. Without loss of generality, we assume that t_i includes $x^i u^i$ as a subword.

We claim that for $i = 1, \dots, n$, $t_i \in T$ if and only if a value of C_i is *true*. This is shown by the induction. For $i = 1, \dots, m$, if $\beta_i = 1$, then t_i must be in T . Thus, by definition of g , $C_i = true$. For $i \geq m + 1$, suppose that for $j, k < i$, this claim is true. We only discuss the case of $t_j \in T$ and $t_k \in T$. Other case is shown in a similar way. By the assumption, values of C_j and C_k is *true*. We remove contributions of t_j and t_k from w_i . The remaining word is $b^j 0 b^k 0 a^i 0 1 dx^i u^i$. Moreover, w_i must has a contribution from L_{bba} . This contribution must be of the form $b^+ 0 b^+ 0 a^+ 1$. Thus, the remaining word after removing this contribution is $0 dx^i u^i$. Therefore, t_i must be in F . On the other hand, a value of $C_i = \text{NOR}(C_j, C_k)$ is *false*. Thus, we hold this claim.

Since t_n must be in T , a value of C_n is *true*. Thus, $g(w) \in CVP$. □

By the discussion above, we can say that \mathcal{L} have a log-space reduction f from CVP and CVP have a log-space reduction g (inverse of f) from \mathcal{L} . □

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