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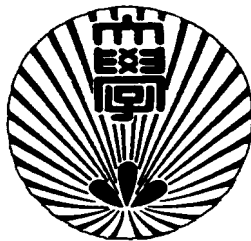
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Abstract

In this paper we deal with 1-D finite cellular automata with a triplet local transition rule 14 and 142. And we investigate their behaviors and they are compared with each other.

1 Introduction

Cellular automata are discrete dynamical systems, which have discrete time and space, and have a local transition rule. Cellular automata were introduced by J. von Neumann as a simulator of a system having self-reproduction and universal computation. Since cellular automata have wide variety and their behavior are similar to those of complex systems as fractal and chaotic phenomenon for the last two decades many researchers investigated their behaviors and applied in mathematics, physics, biology, computer science and so on[1, 2, 3, 4]. Although Cellular automata have simple structure, their behaviors are very complicated. Behaviors of linear cellular automata were investigated by using algebraic methods[5, 8, 9, 10]. But those of many nonlinear cellular automata except some cellular automata[6, 7] are not analysed yet since we cannot investigate them by algebraic methods.

The transition rule of nonlinear cellular automata is difficult to deal with because of not to be able to use algebraic methods. But their behaviors are not always difficult. The numbers of limit cycles and transient length of many nonlinear cellular automata obey simple rules in particular.

In this paper behaviors of the nonlinear cellular automata $CA-14_{\alpha-\beta}(m)$ and $CA-142_{\alpha-\beta}(m)$ are compared while being investigated. We investigate their behaviors and show rules of the number of limit cycles, period length and transient length.

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2 Preliminaries

In this section, we define 1-dimensional finite cellular automata and necessary notations for after discussion.

X is the set $\{0, 1\}^m$, which is called a configuration space, and an element in X is called a configuration. Usually a configuration (x_1, x_2, \dots, x_m) denotes $x_1x_2 \cdots x_m$ for short.

The triplet local transition rule f is a function $\{0, 1\}^3 \longrightarrow \{0, 1\}$, and f is represented as follows:

$$\begin{array}{c} \left| \begin{array}{cccccccc} 111 & 110 & 101 & 100 & 011 & 010 & 001 & 000 \\ r_7 & r_6 & r_5 & r_4 & r_3 & r_2 & r_1 & r_0 \end{array} \right| \end{array}$$

Rule number R of f is defined as follows:

$$R = 2^7 r_7 + 2^6 r_6 + \cdots + 2^0 r_0.$$

For instance,

$$14 = 2^3 \times 0 + 2^2 \times 1 + 2^1 \times 1$$

and

$$142 = 2^7 \times 0 + 2^3 \times 1 + 2^2 \times 1 + 2^1 \times 1,$$

so the triplet local transition rule of rule number 14 and 142 are illustrated as follows:

$$\begin{array}{l} \text{Rule 14} \\ \text{Rule 142} \end{array} \left| \begin{array}{cccccccc} 111 & 110 & 101 & 100 & 011 & 010 & 001 & 000 \\ 0 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \\ 1 & 0 & 0 & 0 & 1 & 1 & 1 & 0 \end{array} \right|$$

A cellular automaton which is dealt with in this paper is a pair (X, δ) where a global transition function $\delta : X \rightarrow X$ is defined by a triplet local transition rule f as follows;

$$\delta(x_1x_2 \cdots x_m) = (f(\alpha x_1x_2)f(x_1x_2x_3) \cdots f(x_{m-1}x_mx_{m+1}))$$

We call a pair (α, β) a boundary. The boundary condition is cyclic if and only if $\alpha = x_m$ and $\beta = x_1$, and the boundary condition is fixed if α and β are fixed in $\{0, 1\}$. And in particular we called the boundary condition $a - b$ if $\alpha = a$ and $\beta = b$. The triplet local transition rule of rule number R (or the global transition function δ defined by f of rule number R) is denoted by “rule R ” for short.

The cellular automaton (X, δ) such that $X = \{0, 1\}^m$ and the rule number of a triplet local transition rule f which define δ is R and its boundary condition is $a - b$ is denoted by $\text{CA-}R_{a-b}(m)$.

The configuration x is on a limit cycle of period length T if there exists a positive integer s such that $\delta^s(x) = x$, and $T = \min\{s \geq 1 | \delta^s(x) = x\}$. And the configurations $x(1), x(2), \dots, x(T-1)$ form a limit cycle of period length T if $x(i+1) = \delta(x(i))$ and $x(T) = x(1)$, and $x(i)$ is on a limit cycle of period length T where $1 \leq i \leq T-1$. A

limit cycle of period length T is denoted by a T -cycle, in particular 1-cycle is called a fixed point. And a number of a limit cycle of period length T is denoted by $\gamma_T(m)$. A configuration x is non-initial if there exists a configuration y such that $x = \delta(y)$. The height $h(x)$ of x from a limit cycle is defined as follows:

$$h(x) = \min\{s \geq 0 \mid \delta^s(x) \text{ is on a limit cycle}\}$$

Then the transient length $H(m)$ of $\text{CA-}R_{a-b}(m)$ is defined as follows;

$$H(m) = \max\{h(x) \mid x \in \{0, 1\}^m\}$$

The symmetric transition rule f^T of a local transition rule f is defined as follows;

$$f^T(x, y, z) = f(z, y, x)$$

The reverse transition rule \bar{f} of a local transition rule f is defined as follows;

$$\bar{f}(x, y, z) = 1 - f(1 - x, 1 - y, 1 - z)$$

For example, the symmetric rule, the reverse rule and the symmetric reverse rule of rule 14 are rule 84, rule 143 and rule 213 respectively. And the symmetric rule, the reverse rule and the symmetric reverse rule of rule 142 are rule 212, rule 142 and rule 212 respectively. They are illustrated as follows;

	111	110	101	100	011	010	001	000
Rule 84	0	1	0	1	0	1	0	0
Rule 143	1	0	0	0	1	1	1	1
Rule 212	1	1	0	1	0	1	0	0
Rule 213	1	1	0	1	0	1	0	1

So $\text{CA-}14_{\alpha-\beta}(m)$, $\text{CA-}84_{\beta-\alpha}(m)$, $\text{CA-}143_{\bar{\alpha}-\bar{\beta}}(m)$ and $\text{CA-}213_{\bar{\beta}-\bar{\alpha}}(m)$ are isomorphic each other, and $\text{CA-}142_{\alpha-\beta}(m)$, $\text{CA-}212_{\beta-\alpha}(m)$, $\text{CA-}142_{\bar{\alpha}-\bar{\beta}}(m)$ and $\text{CA-}212_{\bar{\beta}-\bar{\alpha}}(m)$ are isomorphic each other.

For the discussion in this paper the following notations are defined; Let A be a subsequence. Then the sequence composed of k A 's, the sequence composed of k bits taken from some A 's and an arbitrary bit are denoted by A^k , A_k^* and $*$, respectively. For example, $(011)^2 = 011011$, $(010)_5^* = 01001$ and $(010)_6^* = (010)^2$.

3 Behaviors of $\text{CA-}14_{\alpha-\beta}(m)$ and $\text{CA-}142_{\alpha-\beta}(m)$

In this section we investigate behaviors of $\text{CA-}14_{\alpha-\beta}(m)$ and $\text{CA-}142_{\alpha-\beta}(m)$ and compare their behaviors with each other. And we investigate the differences between their behaviors caused by the difference between rule 14 and rule 142.

The number of limit cycles and transient length of $\text{CA-}14_{\alpha-\beta}(m)$ and $\text{CA-}142_{\alpha-\beta}(m)$ are as table 1 and table 2.

$\alpha-\beta$	m	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0-0	1 - cycle	2	2	3	3	4	4	5	5	6	6	7	7	8	8	9
	tran.len.	0	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0-1	1 - cycle	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
	tran.len.	1	3	5	7	9	11	13	15	17	19	21	23	25	27	29
1-0	1 - cycle	1	2	2	3	3	4	4	5	5	6	6	7	7	8	8
	tran.len.	1	1	3	4	5	6	7	8	9	10	11	12	13	14	15
1-1	1 - cycle	1	1	1	1	1	1	1	1	1	1	1	1	1	1	1
	tran.len.	1	2	4	6	8	10	12	14	16	18	20	22	24	26	28

Table 1: CA-14 $_{\alpha-\beta}(m)$

$\alpha-\beta$	m	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0-0	1 - cycle	2	2	3	3	4	4	5	5	6	6	7	7	8	8	9
	tran.len.	0	2	3	4	5	6	7	8	9	10	11	12	13	14	15
0-1	1 - cycle	1	2	2	3	3	4	4	5	5	6	6	7	7	8	8
	tran.len.	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1-0	1 - cycle	1	2	2	3	3	4	4	5	5	6	6	7	7	8	8
	tran.len.	1	2	3	4	5	6	7	8	9	10	11	12	13	14	15
1-1	1 - cycle	2	2	3	3	4	4	5	5	6	6	7	7	8	8	9
	tran.len.	0	2	3	4	5	6	7	8	9	10	11	12	13	14	15

Table 2: CA-142 $_{\alpha-\beta}(m)$

From table 1 and table 2 we can easily see that CA-14 $_{\alpha-\beta}(m)$ and CA-142 $_{\alpha-\beta}(m)$ behave regularly. And CA-14 $_{\alpha-0}(m)$ and CA-142 $_{\alpha-0}(m)$ behave the same but CA-14 $_{\alpha-1}(m)$ and CA-142 $_{\alpha-1}(m)$ do differently. So it is conjectured that at boundary condition $\alpha - 0$ there exist no effects on the number of limit cycles and transient length of CA-14 $_{\alpha-0}(m)$ and CA-142 $_{\alpha-0}(m)$ by the different between rule 14 and rule 142. In order to show it we will investigate their behaviors in detail.

First necessary and sufficient conditions for a configuration c to be a fixed point are as follows;

Lemma 1 1. The configuration c is a fixed point of CA-14 $_{\alpha-\beta}(m)$ if and only if c satisfies the following conditions.

- $\alpha c \beta$ contains no 001.
- $\alpha c \beta$ contains no 110.
- $\alpha c \beta$ contains no 111.

2. The configuration c is a fixed point of CA-142 $_{\alpha-\beta}(m)$ if and only if c satisfies the following conditions;

- $\alpha c\beta$ contains no 001.
- $\alpha c\beta$ contains no 110.

Proof.

1. Let $d = \delta(c)$.
 (sufficient condition) Since $f(111) = f(110) = 0$ and $f(001) = 1$ if $\alpha c\beta$ contains 111, 110 or 001, then c is not a fixed point. So if c is a fixed point then $\alpha c\beta$ contains no 111, no 110 and no 001.
 (necessary condition) Note that $f(xyz) = y$ for $xyz \neq 111, 110$ and 001. So if $\alpha c\beta$ contains no 111, no 110 and no 001, then trivially c is a fixed point.
2. (sufficient condition) Assume that $\alpha c\beta$ contains 001 or 110. Since $f(110) = 0$ and $f(001) = 1$ c is not a fixed point.
 (necessary condition) Assume that $\alpha c\beta$ contains no 001 and no 110. Then c is a fixed point since $f(xyz) = y$ for $xyz \neq 001$ and 110.

□

By the above lemma we can easily get the configurations which are fixed points of CA-14 $_{\alpha-\beta}(m)$ and CA-142 $_{\alpha-\beta}(m)$ as the following corollaries.

Corollary 1 1. All fixed points of CA-14 $_{0-0}(m)$ are $(10)^i 0^{m-2i}$ and $(10)_m^*$ where $0 \leq i \leq \lfloor \frac{m-1}{2} \rfloor$.

2. All fixed points of CA-14 $_{1-0}(m)$ are $(01)^i 0^{m-2i}$ where $0 \leq i \leq \lfloor \frac{m}{2} \rfloor$.

3. The configuration $(10)_m^*$ is the unique fixed point of CA-14 $_{0-1}(m)$.

4. The configuration $(01)_m^*$ is the unique fixed point of CA-14 $_{1-1}(m)$.

Corollary 2 1. All fixed points of CA-142 $_{0-0}(m)$ are $(10)^i 0^{m-2i}$ and $(10)_m^*$ where $0 \leq i \leq \lfloor \frac{m}{2} \rfloor$.

2. All fixed points of CA-142 $_{1-0}(m)$ are $(01)^i 0^{m-2i}$ where $0 \leq i \leq \lfloor \frac{m}{2} \rfloor$.

3. All fixed points of CA-142 $_{0-1}(m)$ are $(10)^i 1^{m-2i}$ where $0 \leq i \leq \lfloor \frac{m}{2} \rfloor$.

4. All fixed points of CA-142 $_{1-1}(m)$ are $(01)^i 1^{m-2i}$ and $(01)_m^*$ where $0 \leq i \leq \lfloor \frac{m}{2} \rfloor$.

Lemma 2 The followings are common behaviors of CA-14 $_{\alpha-0}(m)$ and CA-142 $_{\alpha-0}(m)$.

1. Let c and c' be configurations such that $c_{i+1} = 0$ and $c' = c_1 c_2 \cdots c_i 0^{m-i}$, $d = \delta(c)$ and $d' = \delta(c')$. Then

$$d_1 d_2 \cdots d_i = d'_1 d'_2 \cdots d'_i.$$

2. Let c be a configuration and $d = \delta(c)$. If $c_{m-i}c_{m-i+1} \cdots c_m = 0^{i+1}$ then $d_{m-i}d_{m-i+1} \cdots d_m = 0^{i+1}$ where $0 \leq i \leq m-1$.
3. Let c be an arbitrary configuration and $d = \delta(c)$. Then the number of the subsequences 01 of $\alpha d 0$ is equal to that of $\alpha c 0$.
4. Let c be an arbitrary configuration and $d = \delta(c)$. Then the number of the subsequences 10 of $\alpha d 0$ is equal to that of $\alpha c 0$.

Proof.

1. It is trivial.
2. It is trivial.
3. First we assume that $c_i c_{i+1} = 01$. Then if $c_{i-1} = 0$ then $d_{i-1} d_i = 01$, and if $c_{i-1} = 1$ then $d_i d_{i+1} = 01$. Next we assume that $d_i d_{i+1} = 01$.

- Case of $\text{CA-14}_{\alpha-0}(m)$
From $d_{i+1} = 1$ we have $c_i c_{i+1} c_{i+2} = 011, 010$ or 001 and from $d_i = 0$ $c_{i-1} c_i c_{i+1} c_{i+2} = 1011, 1010$ or $*001$.
- Case of $\text{CA-142}_{\alpha-0}(m)$
From $d_{i+1} = 1$ we have $c_i c_{i+1} c_{i+2} = 111, 011, 010$ or 001 . $c_i c_{i+1} c_{i+2} = 111$ contradicts with $d_i = 0$. If $d_i d_{i+1} = 01$ then $c_{i-1} c_i c_{i+1} c_{i+2} = 1011, 1010$ or $*001$.

So the subsequence 01 is derived from itself and results to itself. Hence the number of 01 of $\alpha d 0$ is equal to that of $\alpha c 0$.

4. First we assume that $c_i c_{i+1} = 10$

- Case of $\text{CA-14}_{\alpha-0}(m)$.
If $c_{i-l} c_{i-l+1} \cdots c_{i-1} = 01^{l-1}$ then $d_{i-l} d_{i-l+1} \cdots d_{i-1} = 110^l$ or 010^l where $1 \leq l \leq i$. If $\alpha c_1 c_2 \cdots c_{i-1} = 1^i$ then $\alpha d_1 d_2 \cdots d_{i+1} = 10^{i+1}$.
- Case of $\text{CA-142}_{\alpha-0}(m)$.
If $c_{i-1} = 1$ then $d_{i-1} d_i d_{i+1} = 100$. And if $c_{i-1} = 0$ then $d_{i-1} d_i d_{i+1} = *10$.

Next we assume that $d_i d_{i+1} = 10$.

- Case of $\text{CA-14}_{\alpha-0}(m)$.
From $d_i = 1$ we have $c_{i-1} c_i c_{i+1} = 011, 010$ or 001 and from $d_{i+1} = 0$ $c_{i-1} c_i \cdots c_{i+j} = 01^j 0$ where $1 \leq j \leq l$ if $d_{i+2} d_{i+3} \cdots d_{i+l+1} = 0^l 1$.
- Case of $\text{CA-142}_{\alpha-0}(m)$.
From $d_i = 1$ we have $c_{i-1} c_i c_{i+1} = 111, 011, 010$ or 001 and from $d_{i+1} = 0$ $c_{i-1} c_i c_{i+1} c_{i+2} = 010*, 0110$ or 1110 .

So the subsequence 10 is derived from itself and results to itself. Hence the number of 10 of $\alpha d0$ is equal to that of $\alpha c0$.

□

We let R be 14 or 142 then the following corollary holds;

Corollary 3 *Let δ and δ' be the transition functions of $CA-R_{\alpha-0}(m)$ and $CA-R_{\alpha-0}(i)$, respectively, c and d be configurations of $CA-R_{\alpha-0}(m)$ such that $c_{i+1}c_{i+2}\cdots c_m = 0^{m-i}$ and $d = \delta(c)$, and c' and d' be configurations of $CA-R_{\alpha-0}(i)$ such that $c' = c_1c_2\cdots c_i$ and $d' = \delta'(c')$. Then*

$$d_1d_2\cdots d_i = d'.$$

And in the following lemma we investigate behaviors of the subsequence 1^l in $CA-14_{\alpha-\beta}(m)$ and $CA-142_{\alpha-0}(m)$.

Lemma 3 1. *For $CA-14_{\alpha-\beta}(m)$ the following hold; Let c be an arbitrary configuration. $\delta(c)$ contains no subsequence 111.*

2. *Let c be a configuration, $d = \delta(c)$ and $c_{i+1}c_{i+2}\cdots c_{i+l}$ and $d_{i'+1}d_{i'+2}\cdots d_{i'+l'}$ be n th subsequences of $\alpha c0$ and $\alpha d0$ such that $c_i c_{i+1}\cdots c_{i+l+1} = 01^l0$ and $d_{i'}d_{i'+1}\cdots d_{i'+l'+1} = 01^{l'}0$, respectively. Then for $CA-142_{\alpha-0}(m)$ if $l = 1$ then $l' = 1$ or 2 , and if $l \geq 2$ then $l' = l - 1$ or l .*

Proof.

1. Let $d = \delta(c)$. From the local transition rule 14 $d_i = 1$ is derived from only $c_{i-1} = 0$. So if $d_{i-1}d_id_{i+1} = 111$ then $c_{i-2}c_{i-1}c_i = 000$. But since $f(000) = 0$ $c_{i-2}c_{i-1}c_i = 000$ contradicts with $d_{i-1}d_id_{i+1} = 111$. Hence d contains no 111.
2. From lemma 2 and their proof we see that the subsequence 01 and 10 are derived from themselves and n th subsequence $d_{i'+1}d_{i'+2}\cdots d_{i'+l'}$ of $\alpha d0$ is derived from n th subsequence $c_{i+1}c_{i+2}\cdots c_{i+l}$ of $\alpha c0$.
 - First we assume that $l = 1$, that is, $c_ic_{i+1}c_{i+2} = 010$. If $c_{i-1} = 1$ or $i = 0$ then $d_id_{i+1}d_{i+2} = 010$, and if $c_{i-1} = 0$ then $d_{i-1}d_id_{i+1}d_{i+2} = 0110$.
 - Next we assume that $l \geq 2$, that is, $c_ic_{i+1}\cdots c_{i+l+1} = 01^l0$. If $c_{i-1} = 1$ or $i = 0$ then $d_id_{i+1}\cdots d_{i+l} = 01^{l-1}0$, and if $c_{i-1} = 0$ then $d_{i-1}d_i\cdots d_{i+l} = 01^l0$.

Hence if $l = 1$ then $l' = 1$ or 2 , and if $l \geq 2$ then $l' = l$ or $l - 1$.

□

The following lemma states how the subsequence 0^l behaves in $CA-14_{\alpha-0}(m)$ and $CA-142_{\alpha-0}(m)$.

- Lemma 4** 1. Let c be a non-initial configuration, $d = \delta(c)$ and $c_{i+1}c_{i+2} \cdots c_{i+l}$ and $d_{i'+1}d_{i'+2} \cdots d_{i'+l'}$ be n th subsequences of $\alpha c0$ and $\alpha d0$ such that $c_i c_{i+1} \cdots c_{i+l+1} = 10^l 1$ and $d_{i'} d_{i'+1} \cdots d_{i'+l'+1} = 10^{l'} 1$, respectively. Then for $CA-14_{\alpha-0}(m)$ if $l = 1$ then $l' = 1$ or 2 , and if $l \leq 2$ then $l' = l - 1$ or l .
2. Let c be a configuration, $d = \delta(c)$ and $c_{i+1}c_{i+2} \cdots c_{i+l}$ and $d_{i'+1}d_{i'+2} \cdots d_{i'+l'}$ be n th subsequences of $\alpha c0$ and $\alpha d0$ such that $c_i c_{i+1} \cdots c_{i+l+1} = 10^l 1$ and $d_{i'} d_{i'+1} \cdots d_{i'+l'+1} = 10^{l'} 1$, respectively. Then for $CA-142_{\alpha-0}(m)$ if $l = 1$ then $l' = 1$ or 2 , and if $l \leq 2$ then $l' = l - 1$ or l .

Proof.

1. From lemma 2 and their proof we see that the subsequence 01 and 10 are derived from themselves and n th subsequence $d_{i'+1}d_{i'+2} \cdots d_{i'+l'}$ of $\alpha d0$ is derived from n th subsequence $c_{i+1}c_{i+2} \cdots c_{i+l}$ of $\alpha c0$.
- First we assume that $l = 1$, that is, $c_i c_{i+1} c_{i+2} = 101$. If $c_{i-1} = 0$ or $i = 0$ then $d_i d_{i+1} d_{i+2} = 101$, and if $c_{i-1} = 1$ then $d_{i-1} d_i d_{i+1} d_{i+2} = 1001$.
 - Next we assume that $l \geq 2$, that is, $c_i c_{i+1} \cdots c_{i+l+1} = 10^l 1$. If $c_{i-1} = 0$ or $i = 0$ then $d_i d_{i+1} \cdots d_{i+l} = 10^{l-1} 1$, and if $c_{i-1} = 1$ then $d_{i-1} d_i \cdots d_{i+l} = 10^l 1$.

Hence if $l = 1$ then $l' = 1$ or 2 , and if $l \geq 2$ then $l' = l$ or $l - 1$.

2. By the same way as 1 we can show it.

□

For only $CA-14_{\alpha-\beta}(m)$ the following lemma holds. The behaviors stated in the following lemma are caused by $f(111) = 0$. So in $CA-142_{\alpha-\beta}(m)$ it does not hold.

Lemma 5 For $CA-14_{\alpha-\beta}(m)$ the following hold;

1. Let c be a configuration and $d = \delta(c)$. If $d_i d_{i+1} = 11$ then $c_{i-1} c_i c_{i+1} = 001$.
2. Let c be a configuration and $e = \delta^2(c)$. If $e_i e_{i+1} e_{i+2} e_{i+3} = 1101$ then $c_{i-2} c_{i-1} c_i c_{i+1} = 1101$ where $2 \leq i \leq m - 3$.
3. Let c be a configurations such that $c_1 c_2 \cdots c_{i+1} = 0^i 1$ ($1 \leq i \leq m - 1$) and $e = \delta^k(c)$ ($i + t \geq 3$). Then $e_1 e_2 \cdots e_{i+k+1}$ contains no subsequence 1101 .

Proof.

1. Let $d_i d_{i+1} = 11$ where $1 \leq i \leq m - 1$. Then $c_{i-1} c_i = 00$ from the local transitions rule 14. So we have $c_{i-1} c_i c_{i+1} = 001$ since $f(000) = 0$ and $f(001) = 1$.

2. Let $d = \delta(c)$. From the local transition rule 14 if $d_i = 1$ then $c_{i-1} = 0$. Now we assume that $e_i e_{i+1} e_{i+2} e_{i+3} = 1101$, then $d_{i-1} = d_i = d_{i+2} = 0$. And since $f(000) = 0$ and $f(010) = 1$ we have $d_{i-1} d_i d_{i+1} d_{i+2} = 0010$. And from $d_{i+1} = 1$ we have $c_i = 0$, and since $f(000) \neq 1$ $c_{i+1} c_{i+2} = 11, 10$ or 01 . But $c_{i+1} c_{i+2} = 01$ contradicts with $d_{i+2} = 0$. So $c_{i+1} c_{i+2} = 11$ or 10 , which don't contradict with $d_{i+2} = 0$. Since $c_{i-1} c_i c_{i+1} \neq 000$ $c_{i-1} = 1$, and similarly $c_{i+2} = 1$. Hence we have $c_{i-2} c_{i-1} c_i c_{i+1} = 1101$ if $e_i e_{i+1} e_{i+2} e_{i+3} = 1101$.
3. We prove it by induction on k . Let c be a configuration such that $c_1 c_2 \cdots c_{i+1} = 0^i 1$, $d = \delta(c)$, $b = \delta^2(c)$, $a = \delta^3(c)$, $e = \delta^k(c)$ and $g = \delta^{k+2}(c)$. First if $i \geq 2$ then $d_1 d_2 \cdots d_{i+2} = 0^{i-1} 110$, so $c_1 c_2 \cdots c_{i+1}$ and $d_1 d_2 \cdots d_{i+2}$ contain no 1101. And if $i = 1$ then $d_1 d_2 d_3 = 110$ and $b_1 b_2 b_3 = 100$. So we have $a_2 = 0$. Hence $b_1 b_2 b_3 b_4$ and $a_1 a_2 \cdots a_5$ contain no 1101. Next we assume that $e_1 e_2 \cdots e_{i+k+1}$ contains no 1101. If $g_1 g_2 \cdots g_{i+k+3}$ contains any 1101 then by 2 $e_1 e_2 \cdots e_{i+k+1}$ contains some 1101. So $g_1 g_2 \cdots g_{i+k+3}$ does not contain any 1101.

□

- Lemma 6** 1. Let c be an arbitrary configuration and $e = \delta^k(c)$. Then for CA-14₀₋₁(m) $e_{m-k} e_{m-k+1} \cdots e_m$ contains no subsequence 000 where $2 \leq k \leq m-1$.
2. Let c be an arbitrary configuration and $d = \delta(c)$. Then for CA-14₀₋₀(m)

$$d_{m-1} d_m \neq 01.$$

3. Let c be a configuration and $d = \delta(c)$. Then for CA-14_{1-\beta}(m)

$$d_1 = 0.$$

Proof.

1. Let $d = \delta(c)$, $b = \delta^2(c)$, $e = \delta^k(c)$ and $g = \delta^{k+1}(c)$. First we assume that $b_{m-2} b_{m-1} b_m = 000$. Then since $d_{m-1} d_m \neq 000$ $d_{m-1} = 1$. Similarly $d_{m-3} d_{m-2} = 11$. But by lemma 3 1 d contains no 111. So we have $b_{m-2} b_{m-1} b_m \neq 000$. Next we assume that $e_{m-k} e_{m-k+1} \cdots e_m$ contains no 000. By the assumption $g_{m-k} g_{m-k+1} g_m$ contains no 000. So we let $g_{m-k-1} g_{m-k} g_{m-k+1} = 000$. If $e_{m-k} e_{m-k+1} e_{m-k+2} \neq 000$ then $e_{m-k-2} e_{m-k-1} e_{m-k} = 111$. But by lemma 3 1 e contains no 111. So $e_{m-k} e_{m-k+1} e_{m-k+2} = 000$. This is contradiction. Hence $g_{m-k-1} g_{m-k} g_{m-k+1} \neq 000$ and $g_{m-k-1} g_{m-k} \cdots g_m$ contains no 000.
2. It is trivial.
3. It is trivial since $f(1 * *) = 0$.

□

By lemma 6 3 we can regard $CA-14_{1-\beta}(m)$ after one step transition as $CA-14_{0-\beta}(m-1)$.

Lemma 7 *Let c be a configuration and $d = \delta(c)$. Then the following hold;*

1. *For $CA-142_{0-0}(m)$ if $c_1c_2 = 01$ then $d_1 = 1$.*
2. *For $CA-142_{1-0}(m)$ if $c_1c_2 = 01$ then $d_1d_2 = 01$.*
3. *For $CA-142_{\alpha-0}(m)$ if $c_1c_2 \cdots c_i = 0^i$ then $d_1d_2 \cdots d_{i-1} = 0^{i-1}$ where $2 \leq i \leq m-1$.*

Proof. It is trivial. □

With respect to transient length the following lemma hold;

Lemma 8 1. *Let $C(m, n) = \{c \mid c \text{ is a non-initial configuration of } CA-14_{0-0}(m) \text{ and } 0c \text{ contains } n \text{ subsequences } 01\}$ and $h'(m, n) = \max\{h(c) \mid c \in C(m, n)\}$. Then for $CA-14_{0-0}(m)$*

$$h'(m, n) \leq m - 1.$$

2. *Let $C(m, n) = \{c \mid c \text{ is a configuration of } CA-142_{\alpha-0}(m) \text{ and } \alpha c0 \text{ contains } n \text{ subsequences } 01\}$ and $h'(m, n) = \max\{h(c) \mid c \in C(m, n)\}$. Then for $CA-142_{\alpha-0}(m)$*

$$h'(m, n) \leq m.$$

3. *Let c be an arbitrary configuration and $e = \delta^{m+k}(c)$ ($k \geq 0$). Then for $CA-14_{0-1}(m)$*

$$e_1e_2 \cdots e_{t+1} = (10)_{k+1}^*.$$

That is, $e_1e_2 \cdots e_{t+1}$ is stable.

Proof.

1. We prove it by induction on n . Let c be an arbitrary configuration in $C(m, n+1)$, and $c_{i-1}c_i$ be n th subsequence 01 of c . By lemma 6 2 if $c_{i+1} = 0$ then the range of i is from $2n-1$ to $m-3$, and if $c_{i+1} = 1$ then it is from $2n-1$ to $m-4$. First we assume that $c_{i+1} = 0$. Then we let $c_{i+j}c_{i+j+1}$ be $(n+1)$ th 01 of c , that is, $c = c_1c_2 \cdots c_{i-1}c_i0^j110^{m-i-j-2}$ or $c = c_1c_2 \cdots c_{i-1}c_i0^j10^{m-i-j-1}$. By lemma 6 2 we have $1 \leq j \leq m-i-2$.

- If $j = 1$ or 2 then there exists a non negative integer $k \leq h'(i+1, n)$ such that $\delta^k(c) = 1(01)^{n-1}00110^{m-2n-3}$ by lemma 2 1, corollary 3 and lemma 4 1. So $\delta^{k+2}(c) = 1(10)^n0^{m-2n-1}$, that is, $\delta^{k+2}(c)$ is a fixed point.
- If $j \geq 3$ then there exists a non negative integer $k \leq h'(i+1, n)$ such that $\delta^k(c) = 1(01)^{n-1}0^l110^{m-2n-l-1}$ where $l \leq j$. So $\delta^{k+l}(c) = 1(01)^n0^{m-2n-1}$.

Next we assume that $c_{i+1} = 1$. Then we let $c_{i+j+1}c_{i+j+2}$ be $(n+1)$ th 01 of c , that is, $c = c_1 \cdots c_{i-1}c_i 10^j 110^{m-i-j-3}$ or $c = c_1 \cdots c_{i-1}c_i 10^j 10^{m-i-j-2}$. By lemma 6 2 we have $1 \leq j \leq m-i-3$.

- If $j = 1$ or 2 then there exists a non negative integer $k \leq h'(i+1, n)$ such that $\delta^k(c) = 1(01)^{n-1}00110^{m-2n-3}$, and $\delta^{k+2}(c)$ is a fixed point.
- If $j \geq 3$ there exists a non negative integer $k \leq h'(i+1, n)$ such that $\delta^k = 1(01)^{n-1}0^l 110^{m-2n-l-1}$ where $l \leq j$. So $\delta^{k+l}(c)$ is a fixed point.

Hence from the above discussion we have

$$h'(m, n+1) \leq \max \left\{ \begin{aligned} &\max_{2n-1 \leq i \leq m-3} \{h'(i+1, n) + 2\}, \\ &\max_{2n-1 \leq i \leq m-3} \{h'(i+1, n) + m - i - 2\}, \\ &\max_{2n-1 \leq i \leq m-4} \{h'(i+1, n) + 2\}, \\ &\max_{2n-1 \leq i \leq m-4} \{h'(i+1, n) + m - i - 3\} \end{aligned} \right\}.$$

So

$$h'(m, n+1) \leq \max_{2n-1 \leq i \leq m-3} \{h'(i+1, n) + 2, h'(i+1, n) + m - i - 2\}.$$

Now it is trivial that $h'(m, 1) = m-2 \leq m-1$ for any m . If $h'(m', n) \leq m'-1$ for any $m' \leq m$ then $h'(i+1, n) + m - i - 2 \leq m-2 \leq m-1$ and $h'(i+1, n) + 2 \leq i+2$. That is, if $h'(m', n) \leq i-1$ for any $m' \leq m$ then we have $h'(m, n+1) \leq m-1$.

2. we prove it by induction on n . Let c be an arbitrary configuration in $C(m, n+1)$, and $c_{i-1}c_i$ be n th subsequence 01 of $0c0$. If $c_{i+1} = 0$ then the range of i is from $2n-1$ to $m-2$, and if $c_{i+1} = 1$ then the range of i is from $2n-1$ to $m-3$. First we assume that $c_{i+1} = 0$, and we let $c_{i+j}c_{i+j+1}$ be $(n+1)$ th 01 of $0c0$. That is, $c = c_1 c_2 \cdots c_i 0^j 1^l 0^{m-i-j-l}$ where $1 \leq j \leq m-i-1$ and $1 \leq l \leq m-i-j$.

- If $j = 1$ and $l = 1$, then there exists a non negative integer $k \leq h'(i, n)$ such that $\delta^k(c) = (10)^n 0^{j'} 1^{l'} 0^{m-2n-j'-l'}$ where $0 \leq j' \leq 1$ and $1 \leq l' \leq 2$. So $\delta^{k+2}(c) = (10)^{n+1} 0^{m-2n-2}$, that is, $\delta^{k+2}(c)$ is a fixed point.
- If $j = 1$ and $l \geq 2$, then there exists a non negative integer $k \leq h'(i, n)$ such that $\delta^k(c) = (10)^n 0^{j'} 1^{l'} 0^{m-2n-j'-l'}$ where $0 \leq j' \leq 1$ and $1 \leq l' \leq l$. So $\delta^{k+j'+l'}(c) = (10)^{n+1} 0^{m-2n-2}$, that is, $\delta^{k+j'+l'}(c)$ is a fixed point.
- If $j \geq 2$ and $l = 1$, then there exists a non negative integer $k \leq h'(i, n)$ such that $\delta^k(c) = (10)^n 0^{j'} 1^{l'} 0^{m-2n-j'-l'}$ where $j' \leq j-1$ and $1 \leq l' \leq 2$. So $\delta^{k+j'+1}(c) = (10)^{n+1} 0^{m-2n-2}$, that is, $\delta^{k+j'+1}(c)$ is a fixed point.
- If $j \geq 2$ and $l \geq 2$, then there exists a non negative integer $k \leq h'(i, n)$ such that $\delta^k(c) = (10)^n 0^{j'} 1^{l'} 0^{m-2n-j'-l'}$ where $0 \leq j' \leq j-1$ and $1 \leq l' \leq l$. So $\delta^{k+j'+l'}(c) = (10)^{n+1} 0^{m-2n-2}$, that is, $\delta^{k+j'+l'}(c)$ is a fixed point.

Next we assume that $c_{i+1} = 1$, and let $c_{i+s+t}c_{i+s+t+1}$ be $(n+1)$ th 01 of 0c0. That is, $c = c_1c_2 \cdots c_i 1^s 0^t 1^l 0^{m-i-s-t-l}$ where $1 \leq s \leq m-i-2$, $1 \leq t \leq m-i-s-1$ and $1 \leq l \leq m-i-s-t$.

- If $t = 1$ and $l = 1$ then there exists a non negative integer $k \leq h'(i + s, n)$ such that $\delta^k(c) = (10)^n 0^{t'} 1^{l'} 0^{m-2n-t'-l'}$ where $t' \leq 1$ and $l' \leq 2$. So $\delta^{k+2}(c) = (10)^{n+1} 0^{m-2n-2}$.
- If $t = 1$ and $l \geq 2$ then there exists a non negative integer $k \leq h'(i + s, n)$ such that $\delta^k(c) = (10)^n 0^{t'} 1^{l'} 0^{m-2n-t'-l'}$ where $t' \leq 1$ and $l' \leq l$. So $\delta^{k+t'+l'}(c) = (10)^{n+1} 0^{m-2n-2}$.
- If $t \geq 2$ and $l = 1$ then there exists a non negative integer $k \leq h'(i + s, n)$ such that $\delta^k(c) = (10)^n 0^{t'} 1^{l'} 0^{m-2n-t'-l'}$ where $t' \leq t-1$ and $l' \leq 2$. So $\delta^{k+t'+1}(c) = (10)^{n+1} 0^{m-2n-2}$.
- If $t \geq 2$ and $l \geq 2$ then there exists a non negative integer $k \leq h'(i + s, n)$ such that $\delta^k(c) = (10)^n 0^{t'} 1^{l'} 0^{m-2n-t'-l'}$ where $t' \leq t-1$ and $l' \leq l$. So $\delta^{k+t'+l'}(c) = (10)^{n+1} 0^{m-2n-2}$.

So by the above discussion we have

$$h'(m, n+1) \leq \max \left\{ \begin{aligned} & \max_{2n-1 \leq i \leq m-2} \{h'(i, n) + 2\}, \\ & \max_{2n-1 \leq i \leq m-2} \{h'(i, n) + m - i\}, \\ & \max_{2n-1 \leq i \leq m-2} \max_{1 \leq s \leq m-i-2} \{h'(i + s, n) + 2\}, \\ & \max_{2n-1 \leq i \leq m-2} \max_{1 \leq s \leq m-i-2} \{h'(i + s, n) + m - i - s\} \end{aligned} \right\}$$

Trivially $h'(m, 1) = m$, and if $h'(m', n) \leq m'$ for $m' \leq m$ then $h'(m, n+1) \leq m$.

3. We prove it by induction on k . Let $d = \delta(c)$, $b = \delta^m(c)$, $e = \delta^{m+k}(c)$ and $g = \delta^{m+k+1}(c)$. First we assume that $c_1 = 1$ then $d_1 = 1$ by $f(01*) = 1$ and we have $b_1 = 1 = (10)_1^*$, and we assume that $c_1c_2 \cdots c_{i+1} = 0^i 1$ ($1 \leq i \leq m$) then $d_1d_2 \cdots d_i = 0^{i-1} 1$ and $a_1 = 1$ where $a = \delta^i(c)$, so we have $b_1 = 1 = (10)_1^*$. Next we assume that $e_1e_2 \cdots e_{k+1} = (10)_{k+1}^*$. If k is even then trivially $g_1g_2 \cdots g_{k+2} = (10)_{k+2}^*$ since $f(1***) = 0$. Otherwise since $e_{k+1}e_{k+2}e_{k+3} \neq 000$ by lemma 6 1 we have $g_{k+2} = 1$. That is, $g_1g_2 \cdots g_{k+2} = (10)_{k+2}^*$.

□

From the above discussion we have the following theorems.

Theorem 1 1. The cellular automaton CA-14₀₋₀(m) has only fixed points and for the number of fixed points the following formula holds;

$$\gamma_1(m) = \left\lceil \frac{m+3}{2} \right\rceil.$$

And its transient length is m .

2. The cellular automaton $CA-14_{0-1}(m)$ has an unique fixed point. And its transient length is $2m - 1$.
3. The cellular automaton $CA-14_{1-0}(m)$ has only fixed points and for the number of fixed points the following formula holds;

$$\gamma_1(m) = \left\lfloor \frac{m+2}{2} \right\rfloor.$$

And its transient length is m .

4. The cellular automaton $CA-14_{1-1}(m)$ has an unique fixed point. And its transient length is $2m - 2$.

Theorem 2 1. The cellular automata $CA-142_{0-0}(m)$ and $CA-142_{1-1}(m)$ have only fixed points and for the number of fixed point the following holds;

$$\gamma_1(m) = \left\lfloor \frac{m+3}{2} \right\rfloor$$

And their transient length is m .

2. The cellular automata $CA-142_{1-0}(m)$ and $CA-142_{0-1}(m)$ have only fixed points and for the number of fixed point the following holds;

$$\gamma_1(m) = \left\lfloor \frac{m+2}{2} \right\rfloor$$

And their transient length is m .

4 Conclusion

In this paper we analysed behaviors of $CA-14_{\alpha-\beta}(m)$ and $CA-142_{\alpha-\beta}(m)$. The difference between rule 14 and rule 142 is the image of the subsequence 111. For the boundary condition $\alpha - 0$ the behaviors after m step transition are not influenced by the difference. But for the boundary condition $\alpha - 1$ since the subsequence 111 does not disappear in $CA-142_{\alpha-1}(m)$ the differences between behaviors of $CA-14_{\alpha-1}(m)$ and $CA-142_{\alpha-1}(m)$ don't disappear.

So at boundary condition $\alpha - 0$ it is not important whether $f(111) = 0$ or $f(111) = 1$ when $f(110) = 0$, $f(101) = 0$, $f(100) = 0$, $f(011) = 1$, $f(010) = 1$, $f(001) = 1$ and $f(000) = 0$. But cellular automata with 1-bit different rule from rule 14 don't behave always like $CA-14_{0-0}(m)$ as table 3

It is guessed that the transition rule of $CA-14_{\alpha-0}(m)$ except for 111 characterize its behaviors. In the future we will analyse the behaviors of other cellular automata in the same point of view and make an universal theory of cellular automata.

	m	1	2	3	4	5	6	7	8	9	10	11	12	13
CA-78 ₀₋₀ (m)	1 - cycle	2	3	4	6	8	11	15	20	27	36	48	64	85
	tran.len.	0	1	3	3	5	5	7	7	9	9	11	11	13
CA-46 ₀₋₀ (m)	1 - cycle	2	2	2	2	2	2	2	2	2	2	2	2	2
	tran.len.	0	2	3	4	5	6	7	8	9	10	11	12	13
CA-30 ₀₋₀ (m)	1 - cycle	2	1	2	1	2	1	2	1	2	1	2	1	2
	2 - cycle	0	1	0	1	0	1	0	1	0	1	0	1	0
	tran.len.	0	1	4	8	14	15	17	24	28	32	34	39	44
CA-6 ₀₋₀ (m)	1 - cycle	2	2	3	3	4	4	5	5	6	6	7	7	8
	tran.len.	0	2	2	6	6	10	10	14	14	18	18	22	22
CA-10 ₀₋₀ (m)	1 - cycle	1	1	1	1	1	1	1	1	1	1	1	1	1
	tran.len.	1	2	3	4	5	6	7	8	9	10	11	12	13
CA-12 ₀₋₀ (m)	1 - cycle	2	3	5	8	13	21	34	55	89	144	233	377	610
	tran.len.	0	1	1	1	1	1	1	1	1	1	1	1	1
CA-15 ₀₋₀ (m)	1 - cycle	1	1	1	1	1	1	1	1	1	1	1	1	1
	tran.len.	1	2	3	4	5	6	7	8	9	10	11	12	13

Table 3: Other cellular automata with 1-bit different rule from rule 14

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