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<http://hdl.handle.net/2324/2998>

出版情報 : DOI Technical Report. 124, 1996-03-16. Department of Informatics, Kyushu University
バージョン :
権利関係 :



Categorical Representation Theorems of Fuzzy Relations

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May 16, 1996

Abstract

This paper provides a notion of Zadeh categories as a categorical structure formed by fuzzy relations with sup-min composition, and proves two representation theorems for Dedekind categories (relation categories) with a unit object analogous to one-point set, and for Zadeh categories without unit objects.

Keywords: fuzzy relations, relation algebras, representation theorem, Dedekind categories, Zadeh category.

1 Introduction

Since Zadeh's invention the concept of fuzzy sets has been extensively investigated in mathematics, science and engineering. The notion of fuzzy relations is also a basic one in processing fuzzy information in relational structures, see e.g. Pedrycz [10]. Goguen [2] generalized the concepts of fuzzy sets and relations taking values on partially ordered sets. Fuzzy relational equations were initiated and applied to medical models of diagnosis by Sanchez [12].

On the other hand theory of relations, namely relational calculus, has a long history, see [8, 13, 14] for more details on the history. Almost modern formalizations of relation algebras are affected by the work of Tarski [15]. Mac Lane [7] and Puppe [11] exposed a categorical basis for calculus of additive relations. Freyd and Scedrov [1] developed and summarized categorical relational calculus, which they called allegories. Concerning with applications to the relational theory of graphs and programs, Schmidt and Ströhlein [13] gave a simple proof of a representation theorem for Boolean relation algebras satisfying Tarski rule and point axiom. Also they wrote an excellent text book [14] on relations and graphs with many useful examples in computer science. Also Kawahara and Mizoguchi [3, 6, 4] developed relational methodology for assertion semantics of programs, theory of graph transformations (or graph grammars) and categorical set theory.

A representation theorem for (homogeneous) fuzzy relation algebras satisfying a point axiom has been given by Kawahara and Furusawa. The aim of the paper is to provide a categorical formalization for fuzzy relations. Fuzzy relations are set-fuctions with values in the unit interval $[0, 1]$, that is, functions $R : X \times Y \rightarrow [0, 1]$. The set of all such fuzzy relations on X constitutes a fuzzy relation category. An algebra of fuzzy relations are naturally equipped with semi-scalar multiplication by scalars in the unit interval. Also, unlike Boolean relation categories, fuzzy relation categories are not Schröder ones. Instead of Schröder categories we

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need to adopt Dedekind categories, named after Jean-Pierre Olivier and Dany Serrato [9], in order to study fuzzy relation categories. This paper is organized as follows:

In section 2 we first recall the fundamentals on relation category, which we will call Dedekind categories following to Jean-Pierre Olivier and Dany Serrato [9], and show a simple representation theorem for Dedekind categories with a unit object satisfying a condition, called the strict point axiom. The section 3 provides a notion and some the basic properties of fuzzy algebras, which are a generalization of algebras formed by set-functions with values in the unit interval $[0, 1]$. In section 4 we first state the definition of Zadeh categories as a categorical structure formed by fuzzy relations with sup-min composition. Then some useful properties on point relations, due to G. Schmidt and T. Ströhlein [13], and a point axiom will be introduced to prove a representation theorem for morphisms in Zadeh categories. In particular, the point axiom induces a function assigning a concrete fuzzy relation between the sets of point relations to an abstract relation in Zadeh categories. In section 5 we show a representation theorem for Zadeh categories satisfying a point axiom without the assumption of existence of unit objects. And it is proved that the representation function is a bijection preserving all operations of Zadeh categories. Therefore the representation function turns out to be a functor. However this functor is generally not an isomorphisms of categories since it maps all objects with no points to the empty set. Finally we prove that the representation theorem coincides with the simple representation theorem in section 2, when the Zadeh category has a unit object satisfying the strict point axiom.

2 Complete Dedekind Categories

In this section we first recall the fundamentals on relation category, which we will call Dedekind categories following to Jean-Pierre Olivier and Dany Serrato [9], and show a simple representation theorem for Dedekind categories with a unit object satisfying the strict point axiom.

Definition 2.1 A complete Dedekind category \mathcal{D} is a category satisfying the following:

- D1. [Complete Distributive Lattice] For all pairs of objects X and Y the homo-set $\mathcal{D}(X, Y)$ consisting of all morphisms of X into Y is a complete distributive lattice with the least morphism O_{XY} and the greatest morphism ∇_{XY} .
- D2. [Involution] An involution $\sharp : \mathcal{D} \rightarrow \mathcal{D}$ is a monotone contravariant functor. That is, for all morphisms $\alpha, \alpha' : X \rightarrow Y, \beta : Y \rightarrow Z$,
 - (a) $(\alpha\beta)^\sharp = \beta^\sharp\alpha^\sharp$, (b) $(\alpha^\sharp)^\sharp = \alpha$, (c) If $\alpha \sqsubseteq \alpha'$, then $\alpha^\sharp \sqsubseteq \alpha'^\sharp$.
- D3. [Dedekind Formula] For all morphisms $\alpha : X \rightarrow Y, \beta : Y \rightarrow Z$ and $\gamma : X \rightarrow Z$ the Dedekind formula $\alpha\beta \sqcap \gamma \sqsubseteq \alpha(\beta \sqcap \alpha^\sharp\gamma)$ holds.
- D4. [Quotient Morphisms] For all morphisms $\beta : Y \rightarrow Z$ and $\gamma : X \rightarrow Z$ a quotient morphism $\gamma \div \beta : X \rightarrow Y$ is a morphism such that $\alpha\beta \sqsubseteq \gamma$ if and only if $\alpha \sqsubseteq \gamma \div \beta$ for all morphisms $\alpha : X \rightarrow Y$. \square

Note that complete distributive lattices are equivalent to complete Brouwerian lattices or complete Heyting algebras.

Throughout the section all discussions will be done in a fixed complete Dedekind category \mathcal{D} . We denote the identity morphism on an object X of \mathcal{D} by id_X . The greatest morphism ∇_{XY} is called the universal morphism and the least morphism O_{XY} the zero morphism. A morphism is nonzero if it is not equal to the zero morphism.

Proposition 2.2 Let $\alpha, \alpha' : X \rightarrow Y$ and $\beta, \beta' : Y \rightarrow Z$ be morphisms in \mathcal{D} .

- (a) $O_{XY}^\sharp = O_{YX}, \nabla_{XY}^\sharp = \nabla_{YX}$ and $\text{id}_X^\sharp = \text{id}_X$.

- (b) $\nabla_{XX}\nabla_{XY} = \nabla_{XY}\nabla_{YY} = \nabla_{XY}$.
- (c) $(\alpha \sqcup \beta)^\sharp = \alpha^\sharp \sqcup \beta^\sharp$ and $(\alpha \sqcap \beta)^\sharp = \alpha^\sharp \sqcap \beta^\sharp$.
- (d) $O_{XY}\beta = O_{XZ}$ and $\alpha O_{YZ} = O_{XZ}$.
- (e) $(\sqcup_\lambda \alpha_\lambda)\beta = \sqcup_\lambda \alpha_\lambda \beta$.
- (f) If $\alpha \sqsubseteq \alpha'$ and $\beta \sqsubseteq \beta'$, then $\alpha\beta \sqsubseteq \alpha'\beta'$.
- (g) If $\alpha \sqcup \alpha' = \nabla_{XY}$, $\alpha \sqcap \alpha' = O_{XY}$ and $\nabla_{XX}\alpha = \alpha$, then $\nabla_{XX}\alpha' = \alpha'$.
- (h) If $\text{id}_X \sqsubseteq \alpha\alpha^\sharp$ and $\alpha^\sharp\alpha' \sqsubseteq \text{id}_Y$, then $\alpha \sqsubseteq \alpha'$ implies $\alpha = \alpha'$.
- (i) If $\alpha^\sharp\alpha \sqsubseteq \text{id}_Y$, then $\alpha(\beta \sqcap \beta') = \alpha\beta \sqcap \alpha\beta'$.
- (j) If $u \sqsubseteq \text{id}_X$ and $v \sqsubseteq \text{id}_X$, then $u^\sharp = uu = u$ and $uv = u \sqcap v$.

Proof. (a) $O_{XY}^\sharp \sqsubseteq O_{YX}^\sharp = O_{YX}$ since $O_{YX} \sqsubseteq O_{XY}^\sharp$, and $\nabla_{YX} = \nabla_{YX}^\sharp \sqsubseteq \nabla_{XY}^\sharp$ since $\nabla_{XY}^\sharp \sqsubseteq \nabla_{YX}$, and $\text{id}_X^\sharp = \text{id}_X^\sharp \text{id}_X = \text{id}_X^\sharp \text{id}_X^\sharp = (\text{id}_X^\sharp \text{id}_X)^\sharp = \text{id}_X^\sharp = \text{id}_X$. (b) Immediate from $\nabla_{XY} = \text{id}_X \nabla_{XY} \sqsubseteq \nabla_{XX} \nabla_{XY}$. (c) First note that $\alpha^\sharp \sqcup \beta^\sharp \sqsubseteq (\alpha \sqcup \beta)^\sharp$. Hence $\alpha \sqcup \beta = \alpha^\sharp \sqcup \beta^\sharp \sqsubseteq (\alpha^\sharp \sqcup \beta^\sharp)^\sharp$ and $(\alpha \sqcup \beta)^\sharp \sqsubseteq (\alpha^\sharp \sqcup \beta^\sharp)^\sharp = \alpha^\sharp \sqcup \beta^\sharp$. (d) It is trivial that $O_{XY} \sqsubseteq O_{XZ} \div \beta$. Hence $O_{XY}\beta \sqsubseteq O_{XZ}$ by Z3. Using this result we have $\alpha O_{YZ} = (O_{YZ}^\sharp \alpha^\sharp)^\sharp = (O_{ZY}^\sharp \alpha^\sharp)^\sharp = O_{ZX}^\sharp = O_{XZ}$. (e) It follows from $(\sqcup_\lambda \alpha_\lambda)\beta \sqsubseteq \gamma \iff \sqcup_\lambda \alpha_\lambda \sqsubseteq \gamma \div \beta \iff \forall \lambda : \alpha_\lambda \sqsubseteq \gamma \div \beta \iff \forall \lambda : \alpha_\lambda \beta \sqsubseteq \gamma \iff \sqcup_\lambda \alpha_\lambda \beta \sqsubseteq \gamma$. (f) If $\beta \sqsubseteq \beta'$, then $\alpha\beta \sqsubseteq \alpha\beta \sqcup \alpha\beta' = \alpha(\beta \sqcup \beta') = \alpha\beta'$ by (d). (g) Note that $\alpha' = \text{id}_X \alpha' \sqsubseteq \nabla_{XX}\alpha'$, and $\nabla_{XX}\alpha' \sqcap \alpha \sqsubseteq \nabla_{XX}(\alpha' \sqcap \nabla_{XX}^\sharp \alpha) = \nabla_{XX}(\alpha' \sqcap \nabla_{XX}\alpha) = \nabla_{XX}(\alpha' \sqcap \alpha) = \nabla_{XX}O_{XY} = O_{XY}$. Hence $\nabla_{XX}\alpha' = \nabla_{XX}\alpha' \sqcap \nabla_{XY} = \nabla_{XX}\alpha' \sqcap (\alpha \sqcup \alpha') = (\nabla_{XX}\alpha' \sqcap \alpha) \sqcup (\nabla_{XX}\alpha' \sqcap \alpha') = \alpha'$. \square

The statement (b) in the last proposition indicates that if $\nabla_{XY} \neq O_{XY}$ then both of X and Y are nonempty.

Proposition 2.3 *Let $\alpha : X \rightarrow Y$ be a morphism such that $\nabla_{XX}\alpha = \alpha$. Then the following three conditions are equivalent: (a) $\text{id}_X \sqsubseteq \alpha\alpha^\sharp$, (b) $\nabla_{XX} = \alpha\alpha^\sharp$, (c) $\nabla_{XX} = \alpha\nabla_{YX}$.*

Proof. (a) \Rightarrow (b) If $\text{id}_X \sqsubseteq \alpha\alpha^\sharp$, then $\nabla_{XX} = \nabla_{XX}\text{id}_X \sqsubseteq \nabla_{XX}\alpha\alpha^\sharp = \alpha\alpha^\sharp$. (b) \Rightarrow (c) If $\nabla_{XX} = \alpha\alpha^\sharp$, then $\nabla_{XX} = \alpha\alpha^\sharp \sqsubseteq \alpha\nabla_{YX}$. (c) \Rightarrow (a) If $\nabla_{XX} = \alpha\nabla_{YX}$, then $\text{id}_X = \text{id}_X \sqcap \nabla_{XX} = \text{id}_X \sqcap \alpha\nabla_{YX} \sqsubseteq \alpha(\alpha^\sharp \text{id}_X \sqcap \nabla_{YX}) = \alpha\alpha^\sharp$. \square

We now introduce three notions of unit objects, I -crisp relations and I -points in Dedekind categories. A unit object is an abstraction of the notion of singleton (or one-point) sets. Two distinct univalent and total relation from a unit object are not mutually disjoint. To this end the notion of I -points are in addition assumed to be I -crisp.

Definition 2.4 Let \mathcal{D} be a complete Dedekind category.

- (a) An object I in \mathcal{D} is called a unit object if $O_{II} \neq \text{id}_I = \nabla_{II}$.
- (b) A morphism $\tau : I \rightarrow X$ is I -crisp if $u\rho \sqsubseteq \tau$ implies $\rho \sqsubseteq \tau$ for a nonzero morphism $u : I \rightarrow I$ and a morphism $\rho : I \rightarrow X$.
- (c) An I -point p of X is an I -crisp morphism $p : I \rightarrow X$ such that $p^\sharp p \sqsubseteq \text{id}_X$, $\text{id}_I = pp^\sharp$.
- (d) A unit object I satisfies the strict point axiom if $\sqcup_{p \in \xi(X)} p^\sharp p = \text{id}_X$ for all objects X , where $\xi(X)$ denotes the set of all I -points of X . \square

Note that an I -point p of X is nonempty from its totality condition $\text{id}_I = pp^\sharp$. Thus if an object has at least one I -point then it is nonempty.

Proposition 2.5 *Let I be a unit object in \mathcal{D} .*

- (a) *If $\tau \sqsubseteq p$ for an I -point $p : I \rightarrow X$ and a morphism $\tau : I \rightarrow X$, then $\tau = wp$ for a unique morphism $w : I \rightarrow I$.*
- (b) *If $p_0 \neq p_1$ for I -points $p_0, p_1 : I \rightarrow X$, then $p_0 \sqcap p_1 = O_{IX}$ and $p_0 p_1^\sharp = O_{II}$.*
- (c) *$(u \sqcap v)\tau = u\tau \sqcap v\tau$ for morphisms $u, v : I \rightarrow I$ and $\tau : I \rightarrow X$.*

Proof. (a) $\tau p^\sharp p \sqsubseteq \tau = \tau \sqcap p \sqsubseteq (\tau p^\sharp \sqcap \text{id}_I)p = \tau p^\sharp p$ by D3. (b) First we show that if $p_0 \sqcap p_1 \neq O_{IX}$ then $p_0 = p_1$. As $p_0 \sqcap p_1 \sqsubseteq p_0$, by (a) there is a unique $w : I \rightarrow I$ such that $p_0 \sqcap p_1 = wp_0$. If $p_0 \sqcap p_1 \neq O_{IX}$, then $w \neq O_{II}$ and so $p_0 \sqsubseteq p_1$, because $wp_0 \sqsubseteq p_1$ and p_1 is I -crisp. Finally if $p_0 \sqcap p_1 = O_{IX}$, then $p_0 p_1^\sharp = p_0 p_1^\sharp \sqcap \text{id}_I \sqsubseteq (p_0 \sqcap \text{id}_I p_1) p_1^\sharp = (p_0 \sqcap p_1) p_1^\sharp = O_{II}$ by D3 and 2.2(a). (c) $(u \sqcap v)\tau \sqsubseteq u\tau \sqcap v\tau \sqsubseteq (u \sqcap v\tau^\sharp)\tau \sqsubseteq (u \sqcap v)\tau$. \square

Set $L = \mathcal{D}(I, I)$ and let $\alpha : X \rightarrow Y$ be a morphism. Then L is a complete distributive lattice by D1. A function $\xi(\alpha) : \xi(X) \times \xi(Y) \rightarrow L$ assigning $\xi(\alpha)(p, q) = p\alpha q^\sharp : I \rightarrow I$ to a pair of I -points $p : I \rightarrow X$ and $q : I \rightarrow Y$, gives an L -fuzzy relation of $\xi(X)$ into $\xi(Y)$ in the sense of [2]. The infimum (the greatest lower bound) and the supremum (the least upper bound) of L -fuzzy relations will be denoted by the symbols \sqcap and \sqcup . We write the category of sets and L -fuzzy relations as \mathbf{Rel}_L .

Theorem 2.6 (*Representation Theorem*) *Let I be a unit object satisfying the strict point axiom. Then every morphism $\alpha : X \rightarrow Y$ has a unique representation*

$$\alpha = \sqcup_{p \in \xi(X)} \sqcup_{q \in \xi(Y)} p^\sharp \xi(\alpha)(p, q) q,$$

where $\xi(\alpha)(p, q) = p\alpha q^\sharp$ for I -points $p : I \rightarrow X$ and $q : I \rightarrow Y$.

Proof. Since $\text{id}_X = \sqcup_{p \in \xi(X)} p^\sharp p$ and $\text{id}_Y = \sqcup_{q \in \xi(Y)} q^\sharp q$ by the strict point axiom we have

$$\begin{aligned} \alpha &= \text{id}_X \alpha \text{id}_Y \\ &= (\sqcup_{p \in \xi(X)} p^\sharp p) \alpha (\sqcup_{q \in \xi(Y)} q^\sharp q) \\ &= \sqcup_{p \in \xi(X)} \sqcup_{q \in \xi(Y)} p^\sharp p \alpha q^\sharp q \\ &= \sqcup_{p \in \xi(X)} \sqcup_{q \in \xi(Y)} p^\sharp \xi(\alpha)(p, q) q. \end{aligned}$$

Finally we show the uniqueness of the representation. Assume $\alpha = \sqcup_{p \in \xi(X)} \sqcup_{q \in \xi(Y)} p^\sharp k_{p,q} q$. Then for all $p_0 \in \xi(X), q_0 \in \xi(Y)$ we have $\xi(\alpha)(p_0, q_0) = p_0 \alpha q_0^\sharp = \sqcup_{p \in \xi(X)} \sqcup_{q \in \xi(Y)} p_0 p^\sharp k_{p,q} q q_0^\sharp = k_{p_0, q_0}$ by 2.5(b). \square

Corollary 2.7 *Let I be a unit object satisfying the strict point axiom.*

- (a) *The function $\xi : \mathcal{D}(X, Y) \rightarrow \mathbf{Rel}_L(\xi(X), \xi(Y))$ is bijective.*
- (b) *$\xi(O_{XY}) = O_{\xi(X)\xi(Y)}$, $\xi(\nabla_{XY}) = \nabla_{\xi(X)\xi(Y)}$ and $\xi(\text{id}_X) = \text{id}_{\xi(X)}$,*
- (c) *$\xi(\alpha \sqcup \alpha') = \xi(\alpha) \cup \xi(\alpha')$ and $\xi(\alpha \sqcap \alpha') = \xi(\alpha) \cap \xi(\alpha')$,*
- (d) *$\xi(\alpha^\sharp) = \xi(\alpha)^\sharp$,*

$$(e) \xi(\alpha\beta) = \xi(\alpha)\xi(\beta).$$

Proof. (b) It is immediate that $\xi(O_{XY})(p, q) = pO_{XY}q^\sharp = O_{II}$. It follows from $\text{id}_I = \text{id}_I \text{id}_I = pp^\sharp qq^\sharp \sqsubseteq p\nabla_{XY}q^\sharp$ that $\xi(\nabla_{XY})(p, q) = \text{id}_I$, and $\xi(\text{id}_X)(p, q) = p\text{id}_Xq^\sharp = pq^\sharp$. (c) $\xi(\alpha \sqcup \alpha')(p, q) = p(\alpha \sqcup \alpha')q^\sharp = p\alpha q^\sharp \sqcup p\alpha'q^\sharp = [\xi(\alpha) \cup \xi(\alpha')](p, q)$ and $\xi(\alpha \sqcap \alpha')(p, q) = \chi(\alpha \sqcap \alpha')q^\sharp = p\alpha q^\sharp \sqcap p\alpha'q^\sharp = [\xi(\alpha) \cap \xi(\alpha')](p, q)$. (d) $\xi(\alpha^\sharp)(q, p) = q\alpha^\sharp p^\sharp = (p\alpha q^\sharp)^\sharp = p\alpha q^\sharp = \xi(\alpha)(p, q) = \xi(\alpha)^\sharp(q, p)$. (e) $\xi(\alpha\beta)(p, r) = p\alpha\beta r^\sharp = \sqcup_{q \in \xi(Y)} p\alpha q^\sharp q\beta r^\sharp = \sqcup_{q \in \xi(Y)} (p\alpha q^\sharp \sqcap q\beta r^\sharp) = [\xi(\alpha)\xi(\beta)](p, r)$ since $p\alpha q^\sharp, q\beta r^\sharp \sqsubseteq \text{id}_I$. \square

From the proof of 2.7(b) it is easy to see that $\nabla_{XY} \neq O_{XY}$ for all nonempty objects X and Y (Cf. a condition P1 in Definition 4.4) if \mathcal{D} has a unit object I satisfying the strict point axiom.

Corollary 2.8 *Let I be a unit object satisfying the strict point axiom. Then the following holds:*

$$(a) \sqcup_{p \in \xi(X)} p = \nabla_{IX},$$

$$(b) \nabla_{XY} = \nabla_{XI} \nabla_{IY},$$

$$(c) \text{ If } Y \text{ has at least one } I\text{-point, then } \nabla_{XY} \nabla_{YZ} = \nabla_{XZ}.$$

Proof. (a) For each I -point $p : I \rightarrow X$ we have $p\nabla_{XI} = \nabla_{II} = \text{id}_I$ by 2.3(c). Hence using the strict point axiom $\nabla_{IX} = \nabla_{IX} \text{id}_X = \sqcup_{p \in \xi(X)} \nabla_{IX} p^\sharp p = \sqcup_{p \in \xi(X)} (p\nabla_{XI})^\sharp p = \sqcup_{p \in \xi(X)} p$. (b) By 2.7(b) we have $\xi(\nabla_{XY})(p, q) = \text{id}_I$ and so

$$\nabla_{XY} = \sqcup_{p \in \xi(X)} \sqcup_{q \in \xi(Y)} p^\sharp q = (\sqcup_{p \in \xi(X)} p)^\sharp (\sqcup_{q \in \xi(Y)} q) = \nabla_{IX}^\sharp \nabla_{IY}$$

by the representation theorem 2.6 and (a). (c) Let q be an I -point of Y . Then $\nabla_{IX} = qq^\sharp \nabla_{IX} \sqsubseteq q\nabla_{YX}$ and $\nabla_{IZ} = qq^\sharp \nabla_{IZ} \sqsubseteq q\nabla_{YZ}$. Hence by using (b) $\nabla_{XZ} = \nabla_{IX}^\sharp \nabla_{IZ} \sqsubseteq (q\nabla_{YX})^\sharp q\nabla_{YZ} = \nabla_{XY} q^\sharp q\nabla_{YZ} \sqsubseteq \nabla_{XY} \nabla_{YZ}$. \square

As a result we have proved that a Dedekind category with a unit object satisfying the strict point axiom is equivalent to a subcategory of a category of L -fuzzy relations.

3 Fuzzy Algebras

This section introduces a notion of fuzzy algebras as a mathematical structure formed by fuzzy sets, and some basic properties of fuzzy algebras are described. Roughly speaking fuzzy algebras are complete distributive lattices with semi-multiplication and semi-Boolean property. Throughout of the paper the unit interval of reals k with $0 \leq k \leq 1$ will be denoted by a symbol $[0, 1]$ and reals $k \in [0, 1]$ will be called scalars.

Definition 3.1 A fuzzy algebra $\mathcal{F} = (\mathcal{F}, \sqsubseteq, \sqcup, \sqcap, \cdot, O, \nabla)$ is an algebraic structure over a nonempty set \mathcal{F} satisfying the following:

F1. [Complete Distributive Lattice] A hexad $(\mathcal{F}, \sqsubseteq, \sqcup, \sqcap, O, \nabla)$ is a complete distributive lattice with the least element O and the greatest element ∇ .

F2. [Semi-Scalar Multiplication] An operation $\cdot : [0, 1] \times \mathcal{F} \rightarrow \mathcal{F}$ is a semi-scalar multiplication on \mathcal{F} . That is,

- (a) $0a = O$ and $1a = a$, (b) $k(k'a) = (kk')a$, (c) $k(\sqcup_\lambda a_\lambda) = \sqcup_\lambda ka_\lambda$ and $k(\sqcap_\lambda a_\lambda) = \sqcap_\lambda ka_\lambda$,
- (d) $(\wedge_\lambda k_\lambda)a = \sqcap_\lambda k_\lambda a$, (e) If $ka \sqsubseteq kb$ and $k > 0$, then $a \sqsubseteq b$.

(The multiplication $k \cdot a$ of $a \in \mathcal{F}$ by a scalar $k \in [0, 1]$ will be written as ka , unless confusion occurs.)

F3. [Semi-Boolean Algebra] If $a \sqcap k\nabla = ka$ for all scalars k , then there is an element b such that $a \sqcup b = \nabla$ and $a \sqcap b = O$. \square

Throughout the rest of the section all discussions will be done in a fixed fuzzy algebra \mathcal{F} .

Proposition 3.2 *Let a, b be elements of a fuzzy algebra \mathcal{F} and k, k' scalars.*

- (a) *If $a \sqsubseteq b$, then $ka \sqsubseteq kb$.*
- (b) *If $k \leq k'$, then $ka \sqsubseteq k'a$. In particular, $ka \sqsubseteq a$ and $kO = O$.*
- (c) *$ka \sqcup k'a = (k \vee k')a$.*
- (d) *If $k\nabla = \nabla$ and $k \neq 1$, then $\nabla = O$.*
- (e) *If $k\nabla = \nabla$ and $\nabla \neq O$, then $k = 1$.*
- (f) *If $k\nabla \sqsubseteq k'\nabla$ and $\nabla \neq O$, then $k \leq k'$.*

Proof. (a) If $a \sqsubseteq b$, then $ka = k(a \sqcap b) = ka \sqcap kb \sqsubseteq kb$ by F2(c). (b) If $k \leq k'$, then $ka = (k \wedge k')a = ka \sqcap k'a \sqsubseteq k'a$ by F2(d). (c) Assume $k \leq k'$. Then $ka \sqcup k'a = k'a = (k \vee k')a$ by (b). (d) Assume $k\nabla = \nabla$ and $0 \leq k < 1$. Then by F2(b) it is trivial that $k^n\nabla = \nabla$ for all natural numbers n . Hence $\nabla = \bigcap_{n \geq 0} k^n\nabla = (\bigwedge_{n \geq 0} k^n)\nabla = 0\nabla = O$ by F2(d) and F2(a). (e) This statement is obviously equivalent to (d). (f) Assume $k\nabla \sqsubseteq k'\nabla$ and $k' < k$. Then $0 \leq k'/k < 1$ and $k\nabla \sqsubseteq k[(k'/k)\nabla]$. Therefore $\nabla \sqsubseteq (k'/k)\nabla$ by F2(e) and so $\nabla = O$ by (d). \square

Following to [5] the concept of crisp elements in fuzzy algebras is defined as follows:

Definition 3.3 An element a of a fuzzy algebra \mathcal{F} is crisp if $a \sqcap k\nabla = ka$ for all scalars $k \in [0, 1]$. \square

In the above definition of crisp elements of fuzzy algebras it is trivial that $ka \sqsubseteq a \sqcap k\nabla$ by 3.2(a) and 3.2(b). Note that $\nabla \sqcap k\nabla = k\nabla$ from $k\nabla \sqsubseteq \nabla$ by F1(b). This means that the universal element ∇ is crisp. Also the zero element O is clearly crisp.

Proposition 3.4 *Let a, b be elements of \mathcal{F} and k a scalar.*

- (a) *If a and b are crisp, then so are $a \sqcup b$ and $a \sqcap b$.*
- (b) *If $a \sqcup b = \nabla$ and $a \sqcap b = O$, then both of a and b are crisp.*
- (c) *If b is crisp and $ka \sqsubseteq b$ for $k > 0$, then $a \sqsubseteq b$.*

Proof. (a) Let a and b be crisp. Then $(a \sqcup b) \sqcap k\nabla = (a \sqcap k\nabla) \sqcup (b \sqcap k\nabla) = ka \sqcup kb = k(a \sqcup b)$ by F1 and F2(c). $(a \sqcap b) \sqcap k\nabla = (a \sqcap k\nabla) \sqcap (b \sqcap k\nabla) = ka \sqcap kb = k(a \sqcap b)$ by F2(c). (b) Assume $a \sqcup b = \nabla$ and $a \sqcap b = O$. First by 3.2(b) $k\alpha \sqsubseteq \alpha$ and $\alpha \sqcap k\beta = O$. Then $\alpha \sqcap k\nabla = \alpha \sqcap k(\alpha \sqcup \beta) = (\alpha \sqcap k\alpha) \sqcup (\alpha \sqcap k\beta) = k\alpha$ by F2(c) and F1. (c) Note that $ka \sqsubseteq k\nabla$ by 3.2(a). As b is crisp we have $ka \sqsubseteq b \sqcap k\nabla = kb$. Hence $a \sqsubseteq b$ by F2(e). \square

From the last proposition 3.4(a) and 3.4(b) it is immediate that the set of all crisp elements in a fuzzy algebra \mathcal{F} forms a Boolean algebra.

4 Zadeh Categories

In this section we first state the definition of Zadeh categories as a categorical structure formed by fuzzy relations with sup-min composition. Then some useful properties on point relations, due to G. Schmidt and T. Ströhlein [13], and a point axiom will be introduced to prove a representation theorem for morphisms in Zadeh categories. In particular, the point axiom induces a function assigning a concrete fuzzy relation between the sets of point relations to an abstract relation in Zadeh categories.

Definition 4.1 A Zadeh category \mathcal{Z} is a complete Dedekind category satisfying the following:
 Z1. [Fuzzy Algebra] For all objects X and Y the homo-set $\mathcal{Z}(X, Y)$ consisting of all morphisms of X into Y is a fuzzy algebra. Its fuzzy algebra structure will be denoted by

$$\mathcal{Z}(X, Y) = (\mathcal{Z}(X, Y), \sqsubseteq, \sqcup, \sqcap, \cdot, O_{XY}, \nabla_{XY}).$$

Z2. [Compatibility with Semi-Scalar Multiplication]

(a) $k(\alpha\beta) = (k\alpha)(k\beta)$, (b) $(k\alpha)^\sharp = k\alpha^\sharp$, (c) $(k\alpha)\beta = (k\alpha)(\beta \sqcap k\nabla_{YZ})$
 for morphisms $\alpha : X \rightarrow Y$, $\beta : Y \rightarrow Z$, and a scalar k . \square

Throughout the rest of the paper all discussions will be done in a fixed Zadeh category \mathcal{Z} .

Proposition 4.2 Let $\alpha : X \rightarrow Y$ and $\beta : Y \rightarrow Z$ be morphisms and k a scalar.

- (a) $(k\alpha)\beta \sqsubseteq k\nabla_{XZ}$ and $\alpha(k\beta) \sqsubseteq k\nabla_{XZ}$.
- (b) If β is crisp, then $(k\alpha)\beta = k(\alpha\beta)$. (If α is crisp, then $\alpha(k\beta) = k(\alpha\beta)$).
- (c) If α and β are crisp, then so are α^\sharp and $\alpha\beta$.
- (d) The identity morphism id_X is crisp.

Proof. (a) From Z2(c), Z2(a) it follows that $(k\alpha)\beta = (k\alpha)(\beta \sqcap k\nabla_{YZ}) \sqsubseteq (k\alpha)(k\nabla_{YZ}) = k(\alpha\nabla_{YZ}) \sqsubseteq k\nabla_{XZ}$. (b) Assume that β is crisp. Then $(k\alpha)\beta = (k\alpha)(\beta \sqcap k\nabla_{YZ}) = (k\alpha)(k\beta) = k(\alpha\beta)$ by Z2(c) and Z2(a). (c) $\alpha^\sharp \sqcap k\nabla_{YX} = (\alpha \sqcap k\nabla_{XY})^\sharp = (k\alpha)^\sharp = k\alpha^\sharp$ by 2.2(a), 2.2(b) and Z2(b). $\alpha\beta \sqcap k\nabla_{XZ} \sqsubseteq \alpha[\beta \sqcap \alpha^\sharp(k\nabla_{XZ})] \sqsubseteq \alpha(\beta \sqcap k\nabla_{YZ}) = \alpha(k\beta) = k(\alpha\beta)$ by D3, 3.2(a), Z2(c) and (b). (d) $\text{id}_X \sqcap k\nabla_{XX} = \text{id}_X \sqcap k(\text{id}_X\nabla_{XX}) = \text{id}_X \sqcap (k\text{id}_X)\nabla_{XX} \sqsubseteq (k\text{id}_X)[(k\text{id}_X)^\sharp\text{id}_X \sqcap \nabla_{XX}] = k\text{id}_X$. \square

In view of [5, 13] the concept of points in Zadeh categories is defined as follows:

Definition 4.3 A point x of an object X is a crisp morphism $x : X \rightarrow X$ such that $x^\sharp x \sqsubseteq \text{id}_X$, $\text{id}_X \sqsubseteq xx^\sharp$ and $\nabla_{XX}x = x$. \square

Note that a point x of X is automatically nonzero from the totality $\text{id}_X \sqsubseteq xx^\sharp$ if X is nonempty. An empty object has no points because points are nonzero by the definition.

By making use of the last definition of points in Zadeh categories we add the following axiom:

Definition 4.4 A Zadeh category \mathcal{Z} satisfies the point axiom if:

- P1. For all objects X and Y the universal morphism ∇_{XY} is nonzero (that is, $\nabla_{XY} \neq O_{XY}$).
- P2. For a nonzero morphism $\alpha : X \rightarrow Y$ there is a scalar $k > 0$ and points $x : X \rightarrow X$ and $y : Y \rightarrow Y$ such that $x\alpha \sqcap \nabla_{XY}y = k\nabla_{XY}y$. \square

In what follows we assume that a fixed Zadeh category \mathcal{Z} satisfies the point axiom above.

Proposition 4.5 *Let $x_0, x_1 : X \rightarrow X$, $y, y_0, y_1 : Y \rightarrow Y$ points, and k and k' scalars. Then the following holds:*

- (a) *If $x_0 \neq x_1$, then $x_0 \sqcap x_1 = O_{XX}$ and $x_0 x_1^\sharp = O_{XX}$.*
- (b) *If $\nabla_{XY} y_0 \sqsubseteq \nabla_{XY} y_1$, then $y_0 = y_1$.*
- (c) *If $x_0^\sharp \nabla_{XY} y_0 \sqsubseteq x_1^\sharp \nabla_{XY} y_1$, then $x_0 = x_1$ and $y_0 = y_1$.*
- (d) *If α is a nonzero crisp morphism and $k\alpha \sqsubseteq k'\alpha'$, then $k \leq k'$.*
- (e) *If $\gamma \sqsubseteq \nabla_{XY} y$ and $\nabla_{XX} \gamma = \gamma$, then there is a unique scalar k such that $\gamma = k \nabla_{XY} y$.*

Proof. (a) First we show that if $x_0 \sqcap x_1 \neq O_{XX}$ then $x_0 = x_1$. Assume that $x_0 \sqcap x_1 \neq O_{XX}$. Then by the point axiom there are a scalar $k > 0$ and points z_0, z_1 such that $z_0(x_0 \sqcap x_1) \sqcap z_1 = kz_1$. Hence $z_1 \sqsubseteq z_0(x_0 \sqcap x_1) \sqsubseteq \nabla_{XX} x_0 \sqcap \nabla_{XX} x_1 = x_0 \sqcap x_1$. Therefore $x_0 = z_1 = x_1$ by 2.2(h). Finally, if $x_0 \sqcap x_1 = O_{XX}$, then $x_0 x_1^\sharp = x_0 x_1^\sharp \sqcap \nabla_{XX} \sqsubseteq (x_0 \sqcap \nabla_{XX} x_1) x_1^\sharp = (x_0 \sqcap x_1) x_1^\sharp = O_{XX}$. (b) Assume that $\nabla_{XY} y \sqsubseteq \nabla_{XY} y_0$. Then $\nabla_{XY} \sqsubseteq \nabla_{XY} y y^\sharp = \nabla_{XY} y_0 y_0^\sharp$ and so $y_0 y_0^\sharp \neq O_{YY}$ from $\nabla_{XY} \neq O_{YY}$. Hence $y = y_0$ by (a). (c) First note that $\nabla_{XY} = \nabla_{XX}^\sharp \nabla_{XY} = (x \nabla_{XX})^\sharp \nabla_{XY} = \nabla_{XX} x^\sharp \nabla_{XY}$ by 2.2(g) and 2.3(c). Assume that $x^\sharp \nabla_{XY} y \sqsubseteq x_0^\sharp \nabla_{XY} y_0$. Then $\nabla_{XY} y = \nabla_{XX} x^\sharp \nabla_{XY} y \sqsubseteq \nabla_{XX} x_0^\sharp \nabla_{XY} y_0 = \nabla_{XY} y_0$. Hence $y = y_0$ by (b). Similarly $x = x_0$. (d) By the point axiom there are a scalar $h > 0$ and points x, y such that $x\alpha \sqcap \nabla_{XY} y = h \nabla_{XY} y$. Then $h \nabla_{XY} = h \nabla_{XY} \nabla_{YY} = h \nabla_{XY} y \nabla_{YY} \sqsubseteq x\alpha \nabla_{YY}$. But $\nabla_{XY} \sqsubseteq x\alpha \nabla_{YY}$ by 3.4(c) since $x\alpha \nabla_{YY}$ is crisp by 4.2(c). Hence $k \nabla_{XY} \sqsubseteq k \nabla_{XY} \alpha \nabla_{YY} = x(k\alpha) \nabla_{YY} \sqsubseteq x(k'\alpha') \sqsubseteq k' \nabla_{XY}$. Therefore $k \leq k'$ by 3.2(f). (e) It is trivial that if $\gamma = O_{XY}$ then $\gamma = 0 \nabla_{XY} y$ by F2(a). Next assume $\gamma \neq 0$. Then by the point axiom there are a scalar $k > 0$ and points x_0, y_0 such that $x_0 \gamma \sqcap \nabla_{XY} y_0 = k \nabla_{XY} y_0$. Note that $\gamma = \nabla_{XX} \gamma = x_0 \nabla_{XX} \gamma = x_0 \gamma$ by 2.3(c). Hence $k \nabla_{XY} y_0 \sqsubseteq x_0 \gamma = \gamma \sqsubseteq \nabla_{XY} y$, which implies $y_0 = y$ by 4.3(b) and so $\gamma = k \nabla_{XY} y$. The uniqueness of k follows from (d). \square

Proposition 4.6 *Let $\chi(X)$ be the set of all points of an object X in \mathcal{Z} . Then the following holds:*

- (a) $\nabla_{XX} = \sqcup_{x \in \chi(X)} x$,
- (b) $\text{id}_X = \sqcup_{x \in \chi(X)} x^\sharp x$.

Proof. (a) Set $\alpha = \sqcup_{x \in \chi(X)} x$. It is clear that $\alpha \sqsubseteq \nabla_{XX}$ and $\nabla_{XX} \alpha = \alpha$. As points are crisp by the definition, α is also crisp by 2.2(d) and by the axiom F3 there is a morphism $\bar{\alpha}$ such that $\alpha \sqcup \bar{\alpha} = \nabla_{XX}$ and $\alpha \sqcap \bar{\alpha} = O_{XX}$. Then $\nabla_{XX} \bar{\alpha} = \bar{\alpha}$ by 3.4(d). Assume $\bar{\alpha} \neq O_{XX}$. By the point axiom P2 there are a scalar $k > 0$ and points $x_0, x_1 : X \rightarrow X$ such that $x_0 \bar{\alpha} \sqcap x_1 = kx_1$. Hence $kx_1 \sqsubseteq x_0 \bar{\alpha} \sqsubseteq \nabla_{XX} \bar{\alpha} = \bar{\alpha}$. But $\bar{\alpha}$ is crisp by 3.4(c) and so $x_1 \sqsubseteq \bar{\alpha}$ by 3.4(e). Therefore $x_1 \sqsubseteq \alpha \sqcap \bar{\alpha} = O_{XX}$, which contradicts to $x_1 \neq O_{XX}$. This proves $\bar{\alpha} = O_{XX}$. (b) It is clear that $\sqcup_{x \in \chi(X)} x^\sharp x \sqsubseteq \text{id}_X$. Using $\nabla_{XX} = \sqcup_{x \in \chi(X)} x$ by (a) we have $\text{id}_X = \text{id}_X \sqcap \nabla_{XX} = \sqcup_{x \in \chi(X)} (\text{id}_X \sqcap x) \sqsubseteq \sqcup_{x \in \chi(X)} (\text{id}_X x^\sharp \sqcap \text{id}_X) x \sqsubseteq \sqcup_{x \in \chi(X)} x^\sharp x$. \square

Thus from the point axiom P2 and 4.5(e) there exists a unique scalar $\chi(\alpha)(x, y)$ such that

$$x\alpha \sqcap \nabla_{XY} y = \chi(\alpha)(x, y) \nabla_{XY} y.$$

Thus $\chi(\alpha)$ defines a fuzzy relation from the set of all points of X into the set of all points of Y . In the next section we will prove that the function $\chi : \mathcal{Z} \rightarrow \mathbf{Rel}_{[0,1]}$ is an embedding of Zadeh categories.

5 Representation Theorem

First we prove the representation theorem for Zadeh categories asserting that every morphism in a Zadeh category satisfying the point axiom, can be represented as a union of pair morphisms with semi-scalar weights.

Theorem 5.1 (*Representation Theorem*) *Every morphism $\alpha : X \rightarrow Y$ in a Zadeh category \mathcal{Z} satisfying the point axiom has a unique representation*

$$\alpha = \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x, y) x^\# \nabla_{XY} y,$$

where $\chi(\alpha)(x, y)$ is a unique scalar such that $x\alpha \sqcap \nabla_{XY} y = \chi(\alpha)(x, y) \nabla_{XY} y$.

Proof. Since $\nabla_{YY} = \sqcup_{y \in \chi(Y)} y$ by 4.6(a) we have

$$\begin{aligned} x\alpha &= x\alpha \sqcap \nabla_{XY} \\ &= x\alpha \sqcap \nabla_{XY} \nabla_{YY} \\ &= x\alpha \sqcap \nabla_{XY} (\sqcup_{y \in \chi(Y)} y) \\ &= \sqcup_{y \in \chi(Y)} (x\alpha \sqcap \nabla_{XY} y) \\ &= \sqcup_{y \in \chi(Y)} \chi(\alpha)(x, y) \nabla_{XY} y \end{aligned}$$

and so

$$\begin{aligned} \alpha &= \text{id}_X \alpha \\ &= \sqcup_{x \in \chi(X)} x^\# x\alpha \\ &= \sqcup_{x \in \chi(X)} x^\# [\sqcup_{y \in \chi(Y)} \chi(\alpha)(x, y) \nabla_{XY} y] \\ &= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} x^\# [\chi(\alpha)(x, y) \nabla_{XY} y] \\ &= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x, y) x^\# \nabla_{XY} y \end{aligned}$$

using $\text{id}_X = \sqcup_{x \in \chi(X)} x^\# x$ by 4.6(b). Finally we show the uniqueness of the representation. Assume $\alpha = \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} k_{x,y} x^\# \nabla_{XY} y$. Then for all $x_0 \in \chi(X), y_0 \in \chi(Y)$ we have

$$x_0 \alpha = \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} k_{x,y} x_0 x^\# \nabla_{XY} y = \sqcup_{y \in \chi(Y)} k_{x_0,y} \nabla_{XY} y$$

and

$$x_0 \alpha \sqcap \nabla_{XY} y_0 = \sqcup_{y \in \chi(Y)} [k_{x_0,y} \nabla_{XY} y \sqcap \nabla_{XY} y_0] = k_{x_0,y_0} \nabla_{XY} y_0.$$

Hence $k_{x,y} = \chi(\alpha)(x, y)$ by 4.5(e). \square

Corollary 5.2 *For all objects X and Y in \mathcal{Z} the function $\chi : \mathcal{Z}(X, Y) \rightarrow \mathbf{Rel}_{[0,1]}(\chi(X), \chi(Y))$ is a bijection.*

Proof. If $\chi(\alpha) = \chi(\alpha')$, then by the last theorem we have

$$\alpha = \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x, y) x^\# \nabla_{XY} y = \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha')(x, y) x^\# \nabla_{XY} y = \alpha',$$

which shows that χ is injective. Given a fuzzy relation $R : \chi(X) \times \chi(Y) \rightarrow [0, 1]$ we set $\alpha_R = \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} R(x, y) x^\# \nabla_{XY} y$. Then by the uniqueness of representation in the last theorem we have $R(x, y) = \chi(\alpha_R)(x, y)$, which means that χ is surjective. \square

Theorem 5.3 *For a morphism α and a set $\{k_\lambda\}_\lambda$ of scalars the identity $\sqcup_\lambda k_\lambda \alpha = (\vee_\lambda k_\lambda) \alpha$ holds.*

Proof. First we show that $\sqcup_{\lambda} k_{\lambda} x^{\#} \nabla_{XY} y = (\vee_{\lambda} k_{\lambda}) x^{\#} \nabla_{XY} y$ for points $x : X \rightarrow X$ and $y : Y \rightarrow Y$. It is obvious that $\sqcup_{\lambda} k_{\lambda} x^{\#} \nabla_{XY} y \sqsubseteq (\vee_{\lambda} k_{\lambda}) x^{\#} \nabla_{XY} y$ and so there is a scalar k such that $\sqcup_{\lambda} k_{\lambda} (x^{\#} \nabla_{XY} y) = k x^{\#} \nabla_{XY} y$. Then $k \leq \vee_{\lambda} k_{\lambda}$ by $k x^{\#} \nabla_{XY} y \sqsubseteq (\vee_{\lambda} k_{\lambda}) x^{\#} \nabla_{XY} y$ and 4.5(b). On the other hand $k_{\lambda} \leq k$ from $k_{\lambda} x^{\#} \nabla_{XY} y \sqsubseteq k x^{\#} \nabla_{XY} y$ and so $\vee_{\lambda} k_{\lambda} \leq k$. Therefore this proves $k = \vee_{\lambda} k_{\lambda}$. We are ready to prove the general case. Since $\alpha = \sqcup_{x,y \in X} \chi(\alpha)(x,y) x^{\#} \nabla_{XY} y$ by the representability theorem 5.1 we have

$$\begin{aligned}
\sqcup_{\lambda} k_{\lambda} \alpha &= \sqcup_{\lambda} k_{\lambda} [\sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x,y) x^{\#} \nabla_{XY} y] \\
&= \sqcup_{\lambda} \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} [k_{\lambda} \chi(\alpha)(x,y) x^{\#} \nabla_{XY} y] \\
&= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \sqcup_{\lambda} [k_{\lambda} \chi(\alpha)(x,y) x^{\#} \nabla_{XY} y] \\
&= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} [(\vee_{\lambda} k_{\lambda}) \chi(\alpha)(x,y) x^{\#} \nabla_{XY} y] \\
&= (\vee_{\lambda} k_{\lambda}) [\sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x,y) x^{\#} \nabla_{XY} y] \\
&= (\vee_{\lambda} k_{\lambda}) \alpha. \square
\end{aligned}$$

The following theorem is known as Tarski rule for Boolean morphism algebras [13, 15].

Theorem 5.4 *For a nonzero morphism $\alpha : X \rightarrow Y$ in \mathcal{Z} there is a scalar $k > 0$ such that $\nabla_{XX} \alpha \nabla_{YY} = k \nabla_{XY}$.*

Proof. By means of the representation theorem 5.1 we have

$$\begin{aligned}
\nabla_{XX} \alpha \nabla_{YY} &= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \nabla_{XX} [\chi(\alpha)(x,y) x^{\#} \nabla_{XY} y] \nabla_{YY} \\
&= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x,y) [\nabla_{XX} x^{\#} \nabla_{XY} y \nabla_{YY}] \\
&= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x,y) \nabla_{XX} \nabla_{XY} \nabla_{YY} \\
&= \sqcup_{x \in \chi(X)} \sqcup_{y \in \chi(Y)} \chi(\alpha)(x,y) \nabla_{XY} \\
&= [\vee_{x \in \chi(X)} \vee_{y \in \chi(Y)} \chi(\alpha)(x,y)] \nabla_{XY}.
\end{aligned}$$

Set $k = \vee_{x \in \chi(X)} \vee_{y \in \chi(Y)} \chi(\alpha)(x,y)$. Then $\nabla_{XX} \alpha \nabla_{YY} = k \nabla_{XY}$. By the way there exist points x, y such that $\chi(\alpha)(x,y) > 0$ by the point axiom P2. Hence $k > 0$, which completes the proof. \square

The following proposition shows that $\chi : \mathcal{Z} \rightarrow \mathbf{Rel}_{[0,1]}$ preserves all operations of fuzzy morphisms, that is, χ is a homomorphism of fuzzy relation algebras.

Proposition 5.5 *Let $\alpha, \alpha' : X \rightarrow Y, \beta : Y \rightarrow Z$ be morphisms and k a scalar. Then the following holds :*

- (a) $\chi(O_{XY}) = O_{\chi(X)\chi(Y)}$, $\chi(\nabla_{XY}) = \nabla_{\chi(X)\chi(Y)}$ and $\chi(\text{id}_X) = \text{id}_{\chi(X)}$.
- (b) $\chi(\alpha \sqcup \alpha') = \chi(\alpha) \cup \chi(\alpha')$.
- (c) $\chi(\alpha \sqcap \alpha') = \chi(\alpha) \cap \chi(\alpha')$.
- (d) $\chi(\alpha^{\#}) = \chi(\alpha)^{\#}$.
- (e) $\chi(k\alpha) = k\chi(\alpha)$.
- (f) $\chi(\alpha\beta) = \chi(\alpha)\chi(\beta)$.

Proof. (a) The first follows from $x O_{XY} \sqcap \nabla_{XY} y = O_{XY} = 0y$, the second follows from $x \nabla_{XY} \sqcap \nabla_{XY} y = \nabla_{XY} \sqcap \nabla_{XY} y = \nabla_{XY} y$ by $x \nabla_{XY} = x \nabla_{XX} \nabla_{XY} = \nabla_{XX} \nabla_{XY} = \nabla_{XY}$, and the last

follows from $x \text{id}_X \sqcap \nabla_{XX} x' = x \sqcap x'$ and 4.5(e).

(b) It follows from

$$\begin{aligned}
\chi(\alpha \sqcup \beta)(x, y) \nabla_{XY} y &= x(\alpha \sqcup \beta) \sqcap \nabla_{XY} y \\
&= (x\alpha \sqcup x\beta) \sqcap \nabla_{XY} y \\
&= (x\alpha \sqcap \nabla_{XY} y) \sqcup (x\beta \sqcap \nabla_{XY} y) \\
&= \chi(\alpha)(x, y) \nabla_{XY} y \sqcup \chi(\beta)(x, y) \nabla_{XY} y \\
&= [\chi(\alpha)(x, y) \vee \chi(\beta)(x, y)] \nabla_{XY} y \\
&= [\chi(\alpha) \cup \chi(\beta)](x, y) \nabla_{XY} y.
\end{aligned}$$

(c) First note that $x(\alpha \sqcap \alpha') = x\alpha \sqcap x\alpha'$. It follows from $x^\sharp x \sqsubseteq \text{id}_X$ that $x(\alpha \sqcap \alpha') \sqsubseteq x\alpha \sqcap x\alpha' \sqsubseteq x(\alpha \sqcap x^\sharp x\alpha') \sqsubseteq x(\alpha \sqcap \alpha')$. Hence we have

$$\begin{aligned}
x(\alpha \sqcap \alpha')(x, y) \nabla_{XY} y &= x(\alpha \sqcap \alpha') \sqcap \nabla_{XY} y \\
&= (x\alpha \sqcap x\alpha') \sqcap \nabla_{XY} y \\
&= (x\alpha \sqcap \nabla_{XY} y) \sqcap (x\alpha' \sqcap \nabla_{XY} y) \\
&= \chi(\alpha)(x, y) \nabla_{XY} y \sqcap \chi(\alpha')(x, y) \nabla_{XY} y \\
&= [\chi(\alpha)(x, y) \wedge \chi(\alpha')(x, y)] \nabla_{XY} y \\
&= [\chi(\alpha) \cap \chi(\alpha')](x, y) \nabla_{XY} y.
\end{aligned}$$

(d) $\alpha \sqcap x^\sharp \nabla_{XY} y = kx^\sharp \nabla_{XY} y$ if and only if $x\alpha \sqcap \nabla_{XY} y = k\nabla_{XY} y$.

First note that $x\alpha \sqcap \nabla_{XY} y = x(\alpha \sqcap x^\sharp \nabla_{XY} y)$ and $\alpha \sqcap x^\sharp \nabla_{XY} y = x^\sharp(x\alpha \sqcap \nabla_{XY} y)$, because $x\alpha \sqcap \nabla_{XY} y \sqsubseteq x(\alpha \sqcap x^\sharp \nabla_{XY} y) \sqsubseteq x x^\sharp(x\alpha \sqcap \nabla_{XY} y) \sqsubseteq \nabla(x\alpha \sqcap \nabla_{XY} y) \sqsubseteq x\alpha \sqcap \nabla_{XY} y$, and $\alpha \sqcap x^\sharp \nabla_{XY} y \sqsubseteq x^\sharp(x\alpha \sqcap \nabla_{XY} y) \sqsubseteq x^\sharp x(\alpha \sqcap x^\sharp \nabla_{XY} y) \sqsubseteq \alpha \sqcap x^\sharp \nabla_{XY} y$ by $x^\sharp x \sqsubseteq \text{id}_X$. Now assume $\alpha \sqcap x^\sharp \nabla_{XY} y = kx^\sharp \nabla_{XY} y$. Then $x\alpha \sqcap \nabla_{XY} y = x(\alpha \sqcap x^\sharp \nabla_{XY} y) = x[k(x^\sharp \nabla_{XY} y)] = k[x(x^\sharp \nabla_{XY} y)] = k\nabla_{XY} y$. Conversely assume $x\alpha \sqcap \nabla_{XY} y = k\nabla_{XY} y$. Then $\alpha \sqcap x^\sharp \nabla_{XY} y = x^\sharp(x\alpha \sqcap \nabla_{XY} y) = x^\sharp(k\nabla_{XY} y) = kx^\sharp \nabla_{XY} y$ (since x^\sharp is crisp).

An identity $\chi(\alpha^\sharp) = \chi(\alpha)^\sharp$ follows from

$$\begin{aligned}
x\alpha \sqcap \nabla_{XY} y = k\nabla_{XY} y &\iff \alpha \sqcap x^\sharp \nabla_{XY} y = kx^\sharp \nabla_{XY} y \\
&\iff \alpha^\sharp \sqcap y^\sharp \nabla_{YX} x = ky^\sharp \nabla_{YX} x \\
&\iff y\alpha^\sharp \sqcap \nabla_{YX} x = k\nabla_{YX} x.
\end{aligned}$$

(e) It follows from

$$\begin{aligned}
\chi(k\alpha)(x, y) \nabla_{XY} y &= x(k\alpha) \sqcap \nabla_{XY} y \\
&= k(x\alpha) \sqcap k\nabla_{XY} y \quad (k(x\alpha) \sqsubseteq k\nabla_{XY} y) \\
&= k(x\alpha) \sqcap k\nabla_{XY} y \\
&= k(x\alpha \sqcap \nabla_{XY} y) \\
&= k[\chi(\alpha)(x, y) \nabla_{XY} y] \\
&= [k\chi(\alpha)(x, y)] \nabla_{XY} y \\
&= [k\chi(\alpha)](x, y) \nabla_{XY} y.
\end{aligned}$$

(f) When $\nabla_{ZZ} = O_{ZZ}$, it is trivial that the equality $\chi(\alpha\beta) = \chi(\alpha)\chi(\beta)$. Therefore we may assume that Z is nonempty. First note that $\chi(\alpha\beta)(x, z) = \vee_y \chi(\alpha y^\sharp y \beta)(x, z)$ from

$$\begin{aligned}
\chi(\alpha\beta)(x, z) \nabla_{XZ} z &= x\alpha\beta \sqcap \nabla_{XZ} z \\
&= x\alpha \text{id}_Y \beta \sqcap \nabla_{XZ} z \\
&= x\alpha(\sqcup_y y^\sharp y)\beta \sqcap \nabla_{XZ} z \\
&= \sqcup_y (x\alpha y^\sharp y \beta \sqcap \nabla_{XZ} z) \\
&= \sqcup_y \chi(\alpha y^\sharp y \beta)(x, z) \nabla_{XZ} z \\
&= [\vee_y \chi(\alpha y^\sharp y \beta)(x, z)] \nabla_{XZ} z. \quad (\text{by 5.3})
\end{aligned}$$

To complete the proof it suffices to show $\chi(\alpha y^\sharp \beta)(x, z) = \chi(\alpha)(x, y) \wedge \chi(\beta)(y, z)$. Note that $x\alpha \sqcap \nabla_{XY}y = x\alpha y^\sharp$ from $x\alpha y^\sharp \sqsubseteq x\alpha \sqcap \nabla_{XY}y \sqsubseteq (x\alpha y^\sharp \sqcap \nabla_{XY}y) = x\alpha y^\sharp$, and $\nabla_{XZ}(z\beta^\sharp \sqcap \nabla_{ZY}y) = \nabla_{XZ}z\beta^\sharp \sqcap \nabla_{XY}y$ from $\nabla_{XZ}(z\beta^\sharp \sqcap \nabla_{ZY}y) \sqsubseteq \nabla_{XZ}z\beta^\sharp \sqcap \nabla_{XZ}\nabla_{ZY}y \sqsubseteq \nabla_{XZ}z\beta^\sharp \sqcap \nabla_{XY}y \sqsubseteq \nabla_{XZ}(z\beta^\sharp \sqcap \nabla_{ZY}y)$. Then we have

$$\begin{aligned} [\chi(\alpha)(x, z) \wedge \chi(\beta)(y, z)]\nabla_{XY}y &= \chi(\alpha)(x, y)\nabla_{XY}y \sqcap \chi(\beta)(y, z)\nabla_{XY}y \quad (\text{by F2(d)}) \\ &= \chi(\alpha)(x, y)\nabla_{XY}y \sqcap \chi(\beta^\sharp)(y, z)\nabla_{XZ}\nabla_{ZY}y \\ &= (x\alpha \sqcap \nabla_{XY}y) \sqcap \nabla_{XZ}(z\beta^\sharp \sqcap \nabla_{ZY}y) \\ &= (x\alpha \sqcap \nabla_{XY}y) \sqcap (\nabla_{XZ}z\beta^\sharp \sqcap \nabla_{XY}y) \\ &= x\alpha y^\sharp \sqcap \nabla_{XZ}z\beta^\sharp. \end{aligned}$$

Set $r = \chi(\alpha y^\sharp \beta)(x, z)$ and $s = \chi(\alpha)(x, y) \wedge \chi(\beta)(y, z)$. Then $x\alpha y^\sharp \beta \sqcap \nabla_{XZ}z = r\nabla_{XZ}z$ and $x\alpha y^\sharp \sqcap \nabla_{XZ}z\beta^\sharp = s\nabla_{XY}y$. We have $r \leq s$ from

$$r\nabla_{XZ}z = x\alpha y^\sharp \beta \sqcap \nabla_{XZ}z \sqsubseteq (x\alpha y^\sharp \sqcap \nabla_{XZ}z\beta^\sharp)\beta \sqsubseteq (s\nabla_{XY}y)\beta \sqsubseteq s\nabla_{XZ},$$

and $s \leq r$ from

$$s\nabla_{XY}y = x\alpha y^\sharp \sqcap \nabla_{XZ}z\beta^\sharp \sqsubseteq (x\alpha z^\sharp \beta \sqcap \nabla_{XZ}z)\beta^\sharp \sqsubseteq (r\nabla_{XZ}z)\beta^\sharp \sqsubseteq r\nabla_{XY}.$$

Therefore $r = s$. \square

The following is immediately deduced from the last proposition.

Corollary 5.6 (*Insertion Theorem*) *Let $\alpha : X \rightarrow Y$ and $\beta : Y \rightarrow Z$ be morphisms and let $x : X \rightarrow X$ and $z : Z \rightarrow Z$ be points. If $\chi(\alpha\beta)(x, z) > 0$, then there is a point $y : Y \rightarrow Y$ such that $\chi(\alpha)(x, y) > 0$ and $\chi(\beta)(y, z) > 0$. \square*

Proposition 5.7 *Let \mathcal{Z} be a Zadeh category satisfying the point axiom and I a unit object in \mathcal{Z} .*

- (a) *If $u : I \rightarrow I$ is a morphism, then there is a scalar k such that $u = k\text{id}_I$.*
- (b) *A morphism $\tau : I \rightarrow X$ is crisp if and only if it is I -crisp.*
- (c) *The unit object I satisfies the strict point axiom.*
- (d) *$\xi(\alpha)(\nabla_{IX}x, \nabla_{IY}y) = \chi(\alpha)(x, y)\text{id}_I$ for I -points x of X and y of Y .*

Proof. (a) Recall that $\text{id}_I = \nabla_{II}$ and a unit object I has just one point id_I . If u is nonzero, then by the point axiom P1 we have $u \sqcap \text{id}_I = k\text{id}_I$ for a scalar $k > 0$. (b) First we show that a crisp morphism $\tau : I \rightarrow X$ is I -crisp. Assume that τ is crisp and $u\tau' \sqsubseteq \tau$ for $\tau' : I \rightarrow X$ and $u \neq \text{id}_I : I \rightarrow I$. By (a) there is a scalar $k > 0$ such that $u = k\text{id}_I$ and so $k\tau' = k(\text{id}_I\tau') = (k\text{id}_I)(k\tau') \sqsubseteq u\tau' \sqsubseteq \tau$. As τ is crisp we have $\tau' \sqsubseteq \tau$ by 3.4(c), which indicates that τ is I -crisp. Conversely we show that an I -crisp morphism is crisp. Assume that $\tau : I \rightarrow X$ is I -crisp. Note that $\tau = \tau\text{id}_X = \sqcup_{x \in \chi(X)} \tau x^\sharp$ by 4.6(b). If $u = \tau x^\sharp$ is nonzero, then $x \sqsubseteq \tau$ because τ is I -crisp, and so $u = \text{id}_I$. Hence $\tau = \sqcup_{x \in \chi(X)} \tau x^\sharp$ and consequently τ is crisp by 3.4(a).

(c) Let $x : X \rightarrow X$ be a point of X and set $p = \nabla_{IX}x$. Then $p^\sharp p = x^\sharp \nabla_{XI} \nabla_{IX}x \sqsubseteq x^\sharp \nabla_{XX}x = x^\sharp x \sqsubseteq \text{id}_X$, $\text{id}_X = \nabla_{II} = \nabla_{IX} \nabla_{XI} \sqsubseteq \nabla_{IX}x x^\sharp \nabla_{XI} = p p^\sharp$ (since X has a point x). As x and ∇_{IX} are crisp, p is crisp by 4.2(c) and so p is I -crisp by (b). This proves that $p = \nabla_{IX}x$ is an I -point of X . Hence $\nabla_{IX} = \nabla_{IX}\text{id}_X = \nabla_{IX}(\sqcup_{x \in \chi(X)} x^\sharp x) = \sqcup_{x \in \chi(X)} \nabla_{IX}x^\sharp x \sqsubseteq \sqcup_{x \in \chi(X)} \nabla_{IX}x \sqsubseteq$

$\sqcup_{p \in \xi(X)} p$, which means $\nabla_{IX} = \sqcup_{p \in \xi(X)} p$.

(d) Set $k = \chi(\alpha)(x, y)$. As has been seen in (c) $\nabla_{IX}x$ and $\nabla_{IY}y$ are I -points and so we have

$$\begin{aligned}
\xi(\alpha)(\nabla_{IX}x, \nabla_{IY}y) &= \nabla_{IX}x\alpha(\nabla_{IY}y)^\sharp \\
&= \nabla_{IX}(x\alpha y^\sharp \sqcap \nabla_{XY})\nabla_{YI} \\
&= \nabla_{IX}(x\alpha y^\sharp \sqcap \nabla_{XY}yy^\sharp)\nabla_{YI} \\
&= \nabla_{IX}(x\alpha \sqcap \nabla_{XY}y)y^\sharp\nabla_{YI} \\
&= \nabla_{IX}(k\nabla_{XY}y)y^\sharp\nabla_{YI} \\
&= k\nabla_{IX}\nabla_{XY}yy^\sharp\nabla_{YI} \\
&= k\nabla_{IX}\nabla_{XY}\nabla_{YI} \\
&= \text{id}_I. \square
\end{aligned}$$

Consequently we have proved that a Zadeh category satisfying the point axiom is equivalent to a subcategory of the category of sets and (ordinary) fuzzy relations. It is an open problem whether a Dedekind category \mathcal{D} with a unit object I satisfying the strict point axiom and with $\mathcal{D}(I, I) \cong [0, 1]$ is a Zadeh category.

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