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## MARKOVIAN DECISION PROCESSES WITH RECURSIVE REWARD FUNCTIONS

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#### 1. Introduction.

In this paper we shall treat Markovian decision processes associated with recursive reward functions, and we shall study the existence of optimal policies in several senses and the properties of them.

This paper clarifies the essence of stochastic dynamic programming problems, and especially exposes the important properties peculiar to a broad class of reward functions which appear in those problems. Such the properties will be introduced, in Section 5, as a recursiveness, a monotonicity and a Lipschitz condition. The recursiveness together with the monotonicity is a basement of Bellman's "principle of optimality" and moreover the Lipschitz condition guarantees the existence of an optimal policy and of the solution to an optimality equation. In Section 5, six examples of generalized reward functions having the recursive property are given, while actually we can conceive infinitely many of them.

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Blackwell [1] formulated the dynamic programming problem of maximizing the expected, discounted sum over the infinite future just in the form of Markovian decision processes, and gave many elegant results. After that, Strauch [2] treated two problems of maximizing the expected, discounted sum and of maximizing the expected, undiscounted but negative sum, in a parallel way. Performance indices, which they wished to maximize in their papers, have the recursive property in our sense (cf. Section 5). The conception of the monotonicity which we shall impose on the generalized reward function is a generalization of the "monotonicity of operators" in their sense. The Lipschitz condition with coefficients whose infinite product converges to zero will play the role similar to the one of "discounting" in their papers. Many of our proofs are owing to their papers, especially to Balckwell [1].

In Section 2, we prepare the probabilistic definitions and notation used throughout the paper. Section 3 is devoted to formulate our decision problems associated with the generalized reward functions. For every  $\varepsilon > 0$ , there always exists a  $(p, \varepsilon)$ -optimal policy (Section 4). In Section 5, we shall define a recursiveness, a monotonicity and a Lipschitz condition, and we shall give an existence theorem for a  $(p, \varepsilon)$ -optimal Markov policy, assuming the recursiveness and Lipschitz condition to the generalized reward function. Examples of recursive reward functions, too, will be given in Section 5. In Section 6 we shall define an operator associated with measurable mapping f from S into A and define an operator associated with a Markov policy. Then it will be shown that there exists a  $(p, \varepsilon)$ -optimal stationary policy under two additional assumptions (the monotonicity and stationarity of g). The relation of the optimal return to the optimality equation will be derived along the same line as in Blackwell [1]. Moreover we shall give the theorem that if the action space A is countable, then for any  $\varepsilon > 0$  there exists an  $\varepsilon$ -optimal stationary policy, and that if A is finite, then there exists an optimal stationary policy.

#### 2. Probabilistic definitions and notation.

In this section we shall give the basic notation and definitions to be used throughout the paper. These follow to those of [1].

By a Borel set we mean a Borel subset of some Polish space. A polish space is the complete separable metric space. Unless otherwise noted, measurable means with respect to the  $\sigma$ -field of Borel subsets of X, which is expressed by  $\mathcal{B}(X)$ . A probability on a non-empty Borel set X is a probability measure defined over the measurable space  $(X, \mathcal{B}(X))$ ; the set of all probabilities on X is denoted by P(X). For any non-empty Borel sets X, Y a conditional probability on Y given X is a function  $q(\cdot|\cdot)$  such that for each  $x \in X$ ,  $q(\cdot|x)$  is a probability on Y and for each  $B \in \mathcal{B}(Y)$ ,  $q(B|\cdot)$  is a Baire function on X. The set of all conditional probabilities on Y given X is denoted by Q(Y|X). The product space of X and Y is denoted by XY. The set of bounded Baire functions on X is denoted by M(X). If  $u, v \in M(X)$ ,  $u \ge v$  means  $u(x) \ge v(x)$  for all  $x \in X$ . For any  $q \in Q(Y|X)$  and q-integrable Baire function on XY, y denotes the Baire function on X whose value at  $x \in X$  is y is y.

 $= \int_Y u(x,y) dq(y\,|\,x). \quad \text{For any } p \in P(X) \text{ and any } u \in M(X), \ pu \text{ is the integral of } u \text{ with respect to } p. \quad \text{For any } p \in P(X), \ q \in Q(Y\,|\,X), \ pq \text{ is the probability on } XY \text{ such that, for every } u \in M(XY), \ pq(u) = p(qu). \quad \text{Every probability } m \text{ on } XY \text{ has a factorization } m = pq \text{ }; \ p \text{ is unique and is just the marginal distribution of the first coordinate variable with respect to } m \text{ }; \ q \text{ is not unique }; \text{ it is a version of the conditional distribution of the second coordinate variable given the first.}$ 

We can extend the above notation in an obvious way to a finite or countable sequence of non-empty Borel sets. The details are ommitted. A  $p \in P(X)$  is degenerate if it is concentrated at some one point  $x \in X$ ; a  $q \in Q(Y|X)$  is degenerate if each  $q(\cdot|x)$  is so. The degenerate q are exactly those for which there is a Baire function f mapping X into Y for which q(f(x)|x)=1 for all  $x \in X$ . Any such f will also denote its associated degenerate q, so that, for any  $u \in M(XY)$ , fu(x)=u(x,f(x)) for all  $x \in X$ .

We shall use the following.

LEMMA 2.1 (Blackwell [1]). For any  $q \in Q(Y|X)$ ,  $\varepsilon > 0$  and q-integrable Baire function u on XY, there is a degenerate  $f \in Q(Y|X)$  such that

$$fu \ge qu$$
 for all  $x \in X$ 

and

$$q(\lbrace y; u(x, y) \geq u(x, f(x)) + \varepsilon \rbrace | x) = 0$$
 for all  $x \in X$ .

#### 3. Decision problem definitions.

A Markovian decision problem is defined by S, A, q, g, where S, A are non-empty Borel sets,  $q \in Q(S|SA)$ , and g a Baire function on  $\bigcup_{n=2}^{\infty} \bigcup_{1 \leq m < n}^{\infty} [S_m \times A_m \times S_{m+1} \times A_{m+1} \times \cdots \times S_n \times R] \cup \bigcup_{m=1}^{\infty} [S_m \times A_m \times S_{m+1} \times A_{m+1} \cdots]]$  where  $S_n = S$ ,  $A_n = A$  for  $n = 1, 2, \cdots$ , and  $R = (-\infty, \infty)$ . A policy  $\pi$  is a sequence  $\{\pi_1, \pi_2, \cdots\}$ , where  $\pi_n \in Q(A|H_n)$  and  $H_n = SA \cdots S$  (2n-1 factors) is the set of possible histories of the system when the n-th actmust be chosen. A policy  $\pi$  is (non-randomized) Markov if each  $\pi_n$  is a degenerate element of Q(A|S), i. e.  $\pi = \{f_1, f_2, \cdots\}$ , where each  $f_n$  is a measurable function from S into A, and is (non-randomized) stationary if there is a measurable function f mapping f into f such that f and f for all f in f the stationary policy defined by f is denoted by f in f and policies f and f and f in f such that f and f in f in f stationary policy defined by f is denoted by f in f in

Any policy  $\pi$ , together with the law of motion q of the system, defines for each initial state s a conditional probability on the set  $ASAS\cdots$  of futures of the system, i.e. it defines an element of  $Q(ASAS\cdots|S)$ , namely  $e_\pi=\pi_1q\pi_2q\cdots$ .  $E^\pi$  and  $E^{(r_m,r_m+1,\ldots,r_n)}$  denote the integral operators by  $e_\pi$  and  $\pi_mq\pi_{m+1}q\cdots\pi_nq$  respectively. Denote the coordinate functions on  $SASA\cdots$  by  $s_1$ ,  $a_1$ ,  $s_2$ ,  $a_2$ ,  $\cdots$  so that our reward on the n-th day, as a function of the past history up to date of the system, is  $g(s_1, a_1, \cdots, a_{n-1}, s_n, v(s_1, a_1, \cdots, a_{n-1}, s_n))$ , and our reward over the infinite future, as a

function on the history  $H = SASAS \cdots$ , is  $g(s_1, a_1, s_2, a_2, \cdots)$ . An expected reward function on S from policy  $\pi$  then is given by

$$I(\pi) = e_{\pi}g = e_{\pi}[g(s_1, a_1, s_2, a_2, \cdots)].$$

For any  $v \in M(H_n)$ , we shall denote by  $I_n(\pi, v)$  the expected reward we could expect if we terminate at the *n*-th stage to receive a generalized terminal reward  $v(h_n) = v(s_1, a_1, \dots, s_n)$  at the terminal state. Thus

$$I_n(\pi, v) = e_{\pi}g(s_1, a_1, \dots, s_n, v(s_1, a_1, \dots, s_n)) = \pi_1 q \pi_2 q \dots \pi_{n-1} q g.$$

We shall denote  $I_n(\pi, 0)$  by  $I_n(\pi)$ . Here it should be noted that  $u \leq v$  may not imply  $I_n(\pi, u) \leq I_n(\pi, v)$  for all  $\pi$  and all n.

For any  $p \in P(S)$ , and any  $\varepsilon > 0$ ,  $\pi^*$  will be called  $(p, \varepsilon)$ -optimal if  $p\{I(\pi^*) \ge I(\pi) - \varepsilon\} = 1$  for every  $\pi$ .  $\pi^*$  will be called  $\varepsilon$ -optimal if it is  $(p, \varepsilon)$ -optimal for every  $p \in P(S)$ , or equivalently, if  $I(\pi^*) \ge I(\pi) - \varepsilon$  for all  $\pi$ , and will be called optimal if it is  $\varepsilon$ -optimal for every  $\varepsilon > 0$ , or equivalently, if  $I(\pi^*) \ge I(\pi)$  for all  $\pi$ .

Throughout this paper we assume  $E^{n\pi}g(h_{n+1})$  is uniformly bounded for all  $\pi$ , n and  $h_{n+1}$ .

#### 4. Existence of a $(p, \varepsilon)$ -optimal policy.

The proof of the following theorem is similar to that of Theorem 1 in [1]. Theorem 4.1. For any  $p \in P(S)$  and any  $\varepsilon > 0$  there is a  $(p, \varepsilon)$ -optimal policy.

Let  $v^* = \sup_{\pi} I(\pi)$ . Then, in general,  $v^*$  may not be measurable. We call the policy  $\pi^*$  strongly  $(p, \varepsilon)$ -optimal if it holds that

$$p\{I(\pi^*) \ge v^* - \varepsilon\} = 1. \tag{4.1}$$

Obviously, the strong  $(p, \varepsilon)$ -optimality implies the  $(p, \varepsilon)$ -optimality. Note that the set  $\{s \; ; \; I(\pi^*) \geq v^* - \varepsilon\}$  is in general not Borel. Strauch [2], however, has shown that it was Borel in the completion of the Borel sets with respect to every  $p \in P(S)$ . Hence the statement (4.1) has the meaning. Furthermore, his method of proof enables us to show

THEOREM 4.2. For any  $p \in P(S)$  and any  $\varepsilon > 0$  there exists a strongly  $(p, \varepsilon)$ -optimal policy.

#### 5. Recursiveness, monotonicity and Lipschitz condition.

Throughout the remainder of this paper, we assume that g satisfies the recursive relations; for any  $\pi$  and  $1 \le m \le n$ 

$$E^{(\pi_m,\pi_{m+1},\cdots,\pi_n,\cdots)}g = E^{(\pi_m,\cdots,\pi_n)}g(s_m, a_m, \cdots, s_{n+1}, E^{n\pi}g).$$
 (5.1)

These relations mean that for  $1 \le m \le n$ ,

$$\pi_m q \pi_{m+1} q \cdots (h_m) g^{(m-1)h}$$

$$= \pi_m q \pi_{m+1} q \cdots \pi_n q(h_m) g(s_m, a_m, \cdots, s_{n+1}, \pi_{n+1} q \pi_{n+2} q \cdots g^{(n)}),$$

where  ${}^{k}h = (s_{k+1}, a_{k+1}, s_{k+2}, \cdots)$  and  $h_{k} = (s_{1}, a_{1}, \cdots, s_{k})$  for  $k = 1, 2, \cdots$ , and  ${}^{0}h = h$ . The

generalized reward function g is called the recursive reward function if g satisfies recursive relations. Note that putting m=1 in (5.1) yields

$$I(\pi) = I_n(\pi, I(^n\pi))$$
.

Before proceeding to  $(p, \varepsilon)$ -domination we shall give six examples of g which is the recursive reward function. The reward functions in the following examples are those which we usually encounter in the dynamic programming problems. (cf. Bellman [4]). The recursiveness in these examples occurs from the peculiarity of the type of g: the accumulation (in some wide sense) of stage-wise rewards.

EXAMPLE 1. (Additive process)

$$g(s_{m}, a_{m}, \dots, s_{n}, v(s_{1}, a_{1}, \dots, s_{n})) = \sum_{k=m}^{n-1} \beta^{k-m} r_{k}(s_{k}, a_{k}, s_{k+1}) + \beta^{n-m} v(s_{1}, a_{1}, \dots, s_{n})$$

$$g(s_{m}, a_{m}, s_{m+1}, a_{m+1}, \dots) = \sum_{k=1}^{\infty} \beta^{k-m} r_{k}(s_{k}, a_{k}, s_{k+1})$$
(5.2)

where  $0 \le \beta \le 1$ ,  $r_k \in M(SAS)$ ,  $v \in M(H_n)$  for  $k = 1, 2, \cdots$ . If  $0 \le \beta < 1$  and if for some  $K < \infty$ ,  $\|r_k\| \le K$   $(k = 1, 2, \cdots)$ , then g in (5.2) is well-defined. The case when  $0 \le \beta < 1$  and  $r_k = r$  for  $k = 1, 2, \cdots$ , consequently g becomes to be well-defined, has been called the discounted case.

EXAMPLE 2. (Absorbed process)

We consider the additive process  $(S, A, q, \{r_k\}, v, \beta = 1)$  where there exist a state  $s_0 \in S$  and a positive  $\alpha < 1$  such that for any action  $a \in A$  and any state  $s \in S$ 

$$q(\lbrace s_0\rbrace \mid s, a) \geq \alpha$$
.

It is convienient to call this process the absorbed process. Let's define new transition law q' as follows:

$$q'(B|s, a) = \begin{cases} \frac{q(B|s, a)}{1-\alpha} & \text{for } s_0 \in B \\ \frac{q(B|s, a) - \alpha}{1-\alpha} & \text{for } s_0 \in B. \end{cases}$$

Note that q' is a Markov transition law. It is easy to show that the absorbed process with  $\{r_k\}$  satisfying  $r_k(s, a, s_0) = 0$  for all  $k \ge 1$  and all  $(s, a) \in SA$  can be reduced to the additive process  $(S, A, q', \{r_k\}, v, \beta')$  with the discount factor  $\beta' = 1 - \alpha < 1$ . This is essentially the result of the fact

$$\int_{S} r_{k}(s_{k}, a_{k}, s_{k+1}) dq(s_{k+1} | s_{k}, a_{k}) = (1-\alpha) \int_{S} r_{k}(s_{k}, a_{k}, s_{k+1}) dq'(s_{k+1} | s_{k}, a_{k}).$$

Furthermore, if  $r_k = r$  for  $k = 1, 2, \cdots$ , then all results established for the discounted Markovian decision processes remain valid for the absorved process, because the absorved process  $(S, A, q, \{r_k\}, v)$  is equivalent to new discounted process  $(S, A, q', \{r_k\}, v, \beta')$ .

EXAMPLE 3. (Multiplicative process)

$$g(s_m, a_m, \dots, s_n, v(s_1, a_1, \dots, s_n)) = \prod_{k=m}^{n-1} r_k(s_k, a_k, s_{k+1}) v(s_1, a_1, s_2, a_2, \dots, s_n),$$

$$g(s_m, a_m, s_{m+1}, a_{m+1}, \dots) = \prod_{k=m}^{\infty} r_k(s_k, a_k, s_{k+1}),$$

where

$$v \in M(H_n)$$
 and  $r_k \in M(SAS)$  for  $k = 1, 2, \dots$ .

We assume herein that  $e_{\pi}g$  is well-defined for every  $\pi$ . We sometimes encounter the multiplicative process in a field of system reliability problems. For example, suppose that the system consists of a series of state, where an arbitrary large number of duplicate components are used. In this system  $r_k(s, a, s') = 1 - (1 - p_k^{n_k(s,a)})$  is the probability of successful performance at the k-th stage, provided that taking an action a at state s and the k-th stage implies we use  $n_k(s, a)$  duplicate components. Here  $p_k$  is the probability of successful operation of the system, when a policy  $\pi$  is used.

EXAMPLE 4. (Multiplicated additive process)

$$\begin{split} g(s_m, \, a_m, \, \cdots, \, s_n, \, v(s_1, \, a_1, \, \cdots, \, s_n)) \\ &= \sum_{j=m}^{n-1} \prod_{k=m}^j r_k(s_k, \, a_k, \, s_{k+1}) + \prod_{k=m}^{n-1} r_k(s_k, \, a_k, \, s_{k+1}) v(s_1, \, a_1, \, \cdots, \, s_n) \,, \\ g(s_m, \, a_m, \, s_{m+1}, \, a_{m+1}, \, \cdots) &= \sum_{j=m}^{\infty} \prod_{k=m}^j r_k(s_k, \, a_k, \, s_{k+1}) \,. \end{split}$$

EXAMPLE 5.

$$\begin{split} g(s_{m},\,a_{m},\,\cdots,\,s_{n},\,v(s_{1},\,a_{1},\,\cdots,\,s_{n})) \\ &= \sum_{j=m}^{n-1} \left\{ \lceil 1 - r_{j}(s_{j},\,a_{j},\,s_{j+1}) \rceil e^{k - m^{\sum_{j=1}^{j} r_{k}(s_{k},\,a_{k},\,s_{k+1})}} \right\} + v(s_{1},\,a_{1},\,\cdots,\,s_{n}) e^{k - m^{\sum_{k=1}^{n-1} r_{k}(s_{k},\,a_{k},\,s_{k+1})}}, \\ g(s_{m},\,a_{m},\,s_{m+1},\,a_{m+1},\,\cdots) &= \sum_{k=m}^{\infty} \left\{ \lceil 1 - r_{j}(s_{j},\,a_{j},\,s_{j+1}) \rceil e^{k - m^{\sum_{j=1}^{j} r_{k}(s_{k},\,a_{k},\,s_{k+1})}} \right\}, \end{split}$$

where

$$v \in M(H_n)$$
 and  $r_n \in M(SAS)$  for  $n = 1, 2, \dots$ 

EXAMPLE 6.

where

$$v \in M(H_n)$$
,  $h \in M(R)$  and  $r_k \in M(SAS)$  for  $k = 1, 2, \dots$ 

It is easy to show that these examples satisfy the recursive relations. We now define a monotonicity and a Lipschitz condition for g.

DEFINITION 5.1. The generalized reward function g is monotone if  $u \le v$   $(u, v \in M(S))$  implies that for any  $\pi$ , all  $h_n$  and all  $n \ge 1$ 

$$\pi_n q(h_n) g(s_n, a_n, s_{n+1}, u(s_{n+1})) \leq \pi_n q(h_n) g(s_n, a_n, s_{n+1}, v(s_{n+1}))$$
.

DEFINITION 5.2. The generalized reward function g satisfies a Lipschitz condition if there exists a sequnce  $\{K_n\}$  of positive numbers such that

$$\begin{split} \|\pi_n q g(s_n, \, a_n, \, s_{n+1}, \, u(s_1, \, a_1, \, \cdots, \, s_{n+1})) \\ & -\pi_n q g(s_n, \, a_n, \, s_{n+1}, \, v(s_1, \, a_1, \, \cdots, \, s_{n+1})) \| \\ & \leq K_n \|u - v\| \end{split}$$

for all  $u, v \in M(H_{n+1})$ , all  $\pi_n \in Q(A|H_n)$  and every  $n \ge 1$ .

Note that the norm  $\|\cdot\|$  at the left hand side in above definition is

$$\sup_{h_n \in H_n} |\pi_n qg(s_n, a_n, s_{n+1}, u(s_1, a_1, \cdots, s_{n+1}))(h_n) - \pi_n qg(s_n, a_n, s_{n+1}, v(s_1, a_1, \cdots, s_{n+1}))(h_n)|.$$

Now we put the following assumptions.

Assumption (I). g is monotone.

Assumption (II). g satisfies the Lipschitz condition with the Lipschitz coefficients such that  $K_1K_2\cdots K_N\to 0$  as  $N\to\infty$ .

THEOREM 5.1. Under Assumption (II), for any  $p \in P(S)$ ,  $\varepsilon > 0$  and any  $\pi$  there exists a Markov policy  $\pi^*$  which  $(p, \varepsilon)$ -dominates  $\pi$ , i. e.  $p\{I(\pi^*) \ge I(\pi) - \varepsilon\} = 1$ .

PROOF. If any two policies  $\pi$ ,  $\pi'$  agree for the first N stages, i.e., if  $\pi = \{\pi_1, \pi_2, \cdots, \pi_N, \pi_{N+1}, \pi_{N+2}, \cdots\}$  and  $\pi' = \{\pi_1, \pi_2, \cdots, \pi_N, \pi'_{N+1}, \pi'_{N+2}, \cdots\}$  then

$$\begin{split} \|I(\pi)-I(\pi')\| &= \|\pi_1 q g(s_1,\ a_1,\ s_2,\ E^{1\pi} g) - \pi_1 q g(s_1,\ a_1,\ s_2,\ E^{1\pi'} g)\| \\ &\leq K_1 \|E^{1\pi} g - E^{1\pi'} g\| \\ &= K_1 \|\pi_2 q g(s_2,\ a_2,\ s_3,\ E^{2\pi} g) - \pi_2 q g(s_2,\ a_2,\ s_3,\ E^{2\pi'} g)\| \\ &\leq K_1 K_2 \|E^{2\pi} g - E^{2\pi'} g\| \\ &\vdots \\ &\leq K_1 K_2 \cdots K_N \|E^{N\pi} g - E^{N\pi'} g\| \leq K_1 K_2 \cdots K_N (\|E^{N\pi} g\| + \|E^{N\pi'} g\|) \\ &\leq 2K_1 K_2 \cdots K_N \cdot L \;. \end{split}$$

where L is a finite number depending on none of  $\pi$ ,  $\pi'$  and N. Hence we may suppose that  $\pi$  is already Markov from some stage on, say for n > N, namely  $\pi = \{\pi_1, \pi_2, \cdots, \pi_N, f_{N+1}, f_{N+2}, \cdots\}$ . By the recursiveness of g,

$$E^{\pi}g = E^{(\pi_1, \dots, \pi_{N-1})}g(s_1, a_1, \dots, s_N, E^{N-1\pi}g)$$
.

On the other hand, by Lemma 2.1, for  $\gamma = \frac{\varepsilon}{N}$  there exists an  $f_N$  such that

$$\begin{split} E^{N-1\pi}g &= \pi_N q g(s_N, \, a_N, \, s_{N+1}, \, E^{N\pi}g(s_{N+1})) \\ & \leq f_N q g(s_N, \, a_N, \, s_{N+1}, \, E^{N\pi}g(s_{N+1})) + \gamma \\ & \text{with } \, p\pi_1 q \cdots n_{N-1} q \text{-prob. } 1 \, . \end{split}$$

-

That is,

$$\pi_{N-1}q\pi_N qg(s_N, a_N, s_{N+1}, E^{N\pi}g) \leq \pi_{N-1}q[f_N qg(s_N, a_N, s_{N+1}, E^{N\pi}g) + \gamma],$$

or equivalently

$$\pi_{N-1}qg(s_{N-1}, a_{N-1}, s_N, E^{N-1\pi}g) \leq \pi_{N-1}qf_Nqg(s_N, a_N, s_{N+1}, E^{N\pi}g) + \gamma$$
 with  $p\pi_1q \cdots \pi_{N-2}q$ -prob. 1.

Using this procedure N times will produce a Markov  $\pi^* = \{f_1, f_2, \dots\}$  such that

$$E^{\pi}g \leq E^{(f_1,f_2,\cdots,f_N,N_{\pi})}g + N\gamma = E^{\pi *}g + \varepsilon$$
 with  $p$ -prob. 1.

This completes the proof.

COROLLARY. Under Assumption (II), for any  $\varepsilon > 0$  there exists a  $(p, \varepsilon)$ -optimal Markov policy  $\pi^*$ .

#### 6. Stationary policies and operators.

We shall say that g is stationary if for every  $m \ge 1$  and  $k \ge 1$ 

$$g \mid S_m \times A_m \times S_{m+1} \times R = g \mid S_{m+k} \times A_{m+k} \times S_{m+k+1} \times R$$

by identifying  $S_m \times A_m \times S_{m+1} \times R$  and  $S_{m+k} \times A_{m+k} \times S_{m+k+1} \times R$ , where  $g \mid E$  denotes the restriction of g on E.

Throughout this section we assume that g is stationary. The stationarity of g yields that  $r_1 = r_2 = \cdots = r_n = \cdots = (\text{say } r)$  in every example stated in Section 5.

DEFINITION 6.1. For each measurable function f mapping S into A, and  $u \in M(S)$ , let

$$T_f u(s) = fqg(s, a, s', u(s'))$$
  
=  $\int_{c} g(s, f(s), s', u(s')) dq(s' | s, f(s)).$ 

It is easily verified that the operator  $T_f$  associated with measurable function f mapping S into A is a mapping from M(S) into M(S).

Now we shall express the operator  $T_f$  explicitly for each example. For Example 1, also for Example 2, we have

$$\begin{split} T_f u(s) &= fq \llbracket r + \beta u \rrbracket(s) \\ &= \int_{\mathcal{S}} \llbracket r(s, f(s), s') + \beta u(s') \rrbracket dq(s' \mid s, f(s)) \; . \end{split}$$

The operator  $T_f$  in this case was originally given by Blackwell [1]. For Example 3, we have

$$T_{f}u(s) = fq[r \cdot u](s)$$

$$= \int_{s} [r(s, f(s), s')u(s')]dq(s'|s, f(s)).$$

For Example 4, we have

$$\begin{split} T_f u(s) &= fq \llbracket r + r \cdot u \rrbracket(s) \\ &= \int_{s} \llbracket r(s, f(s), s') + r(s, f(s), s') u(s') \rrbracket dq(s' \mid s, f(s)) \,. \end{split}$$

For Example 5, we have

$$\begin{split} T_f u(s) &= fq [(1-r)e^r + e^r u](s) \\ &= \int_S [(1-r(s,\,f(s),\,s'))e^{r(s,f(s),s')} \\ &\quad + e^{r(s,f(s),s')}u(s')]dq(s'\,|\,s,\,f(s)) \,. \end{split}$$

And for Example 6 we have

$$\begin{split} T_f u(s) &= fq \Big[ h(r) + \frac{u}{r} \Big](s) \\ &= \int_{s} \Big[ h(r(s, f(s), s')) + \frac{u(s')}{r(s, f(s), s')} \Big] dq(s' \mid s, f(s)) \; . \end{split}$$

The following two lemmas are immediate consequences of the assumptions. Lemma 6.1. Under Assumption (I), the operator  $T_f$  is monotone for any measurable f from S into A. That is,  $u \leq v$  implies

$$T_f u \leq T_f v$$
.

LEMMA 6.2. Under Assumption (II),  $T_f$  is a contraction mapping on M(S) with some contraction coefficient 0 < K < 1 for all measurable f from S into A.

We now return to Examples and look over whether the operator  $\mathcal{T}_f$  is a monotone or contraction mapping.

In Example 1,  $T_f$  is always a monotone and contraction mapping on M(S) for  $0 \le \beta < 1$ . If  $r \ge 0$  then,  $T_f$  in both of Example 3 and Example 4 is monotone.  $T_f$  is always monotone for Example 5. If  $\|r\| < 1$ , then  $T_f$  in both of Example 3 and Example 4 is a contraction mapping. In Example 5, if  $\|e^r\| < 1$  then  $T_f$  is a contraction. In Example 6, if r > 0 then  $T_f$  is monotone, and if  $\inf_{(s,a,s') \in SAS} |r(s,a,s')| > 1$  then  $T_f$  is a contraction.

Theorem 6.1. (a) Under Assumption (II) for any  $u \in M(S)$  and constant c,  $\|T_f(u+c)-T_fu\| \leq Kc$ .

(b) For any Markov policy  $\pi = \{f_1, f_2, \dots\}$ ,  $TI(\pi) = I(f, \pi)$ , where  $\{f, \pi\}$  denotes the Markov one  $\{f, f_1, f_2, \dots\}$ .

For any Markov  $\pi = \{f_1, f_2, \cdots\}$  we shall say that f mapping S into A is  $\pi$ -generated if there exists a partition of S into Borel sets  $S_1, S_2, \cdots$ , such that  $f = f_n$  on  $S_n$ ; we say that a Markov  $\pi' = \{g_1, g_2, \cdots\}$  is  $\pi$ -generated if each  $g_n$  is so. We associate with each Markov  $\pi = \{f_1, f_2, \cdots\}$  the operator  $U_{\pi}$ , mapping M(S) into M(S), defined by  $U_{\pi}u = \sup_{n \ge 1} T_{f_n}u$ .

Theorem 6.2. (a) If  $T_{f_n}$  is monotone for  $n \ge 1$ , then  $U_{\pi}$  is so.

(b) If 
$$||T_{f_n}u - T_{f_n}v|| \le K||u - v||$$
 for  $k = 1, 2, \dots$ , then  $||U_{\pi}u - U_{\pi}v|| \le K||u - v||$ .

- (c) For any  $T_f$  assoicated with  $\pi$ -generated f,  $T_f u \leq U_{\pi} u$ .
- (d) For any  $u \in M(S)$  and any  $\varepsilon > 0$ , there exists a  $\pi$ -generated f whose associated  $T_f$  satisfies  $T_f u \ge U_{\pi} u \varepsilon$ .

PROOF. (a) is trivial. (c) is immediate from the definition. For (b), consider arbitrary u, v and s, and write  $(U_{\pi}u)(s) = (U_{\pi}v)(s) + k$ . Consider the case k > 0. For any integer n, there exists an m such that

$$T_{f_m}u(s) \ge U_\pi u(s) - \frac{k}{n} . \tag{6.1}$$

On the other hand,

$$U_{\pi}u(s) - \frac{k}{n} \ge U_{\pi}u(s) - k = U_{\pi}v(s) \ge T_{fm}v(s)$$
. (6.2)

Hence, (6.1), (6.2) and the assumption yield

$$0 \leq U_{\pi}u(s) - U_{\pi}v(s) - \frac{k}{n} \leq T_{f_m}u(s) - T_{f_m}v(s)$$

$$\leq K \|u - v\|. \tag{6.3}$$

Since above (6.3) is valid for each n, it follows that

$$|U_{\pi}u(s)-U_{\pi}v(s)| \leq K||u-v||$$
.

This inequality is trivial for k=0 and is similarly established for k<0. Then

$$||U_{\pi}u - U_{\pi}v|| \leq K||u - v||$$
.

For (d), the proof is similar to Theorem 4 (d) in [1].

THEOREM 6.3. (a) For any  $u \in M(S)$  and  $\varepsilon > 0$ , there exists a  $\pi$ -generated f whose associated  $T_f$  satisfies  $||T_f u - U_\pi u|| \le \varepsilon$ .

(b) For any  $u \in M(S)$ ,  $n \ge 1$  and  $\varepsilon > 0$ , there exists a  $\pi$ -generated  $\pi'$  such that  $||I_n(\pi', u) - U_n^n u|| \le \varepsilon$ .

PROOF. (a) is immediate from Theorem 6.2 (d). For (b), then there exists a  $\pi$ -generated  $f_n$  such that

$$||T_{f'_n}u-U_{\pi}u|| \leq \varepsilon_n$$
.

Similarly for  $\varepsilon_{n-1} > 0$  there exists  $\pi$ -generated  $f'_{n-1}$  such that

$$||T_{f'_{n-1}}(T_{f'_n}u)-U_{\pi}(T_{f'_n}u)|| \leq \varepsilon_{n-1}.$$

Then

$$\begin{split} \|T_{f'_{n-1}}T_{f'_{n}}u - U_{\pi}^{2}u\| & \leq \|T_{f'_{n-1}}T_{f'_{n}}u - U_{\pi}(T_{f'_{n}}u)\| + \|U_{\pi}(T_{f'_{n}}u) - U_{\pi}(U_{\pi}u)\| \\ & \leq \varepsilon_{n-1} + K\|T_{f'_{n}}u - U_{\pi}u\| \\ & \leq \varepsilon_{n-1} + K\varepsilon_{n} \,. \end{split}$$

In general there exist  $f_1'$ ,  $f_2'$ ,  $\cdots$ ,  $f_n'$  such that for  $i=1, 2, \cdots$ , n

$$||T_{f_1'} \cdots T_{f_n'} u - U_{\pi}^{n-i+1} u|| \leq d_i$$
,

where  $d_i = \varepsilon_i + K\varepsilon_{i+1} + \cdots + K^{n-i}\varepsilon_n$ . In particular we can choose  $\varepsilon_1, \cdots, \varepsilon_n$  so that  $d_1 \le \varepsilon$ . This showes the existence of the desired  $\pi'$ .

Note that Theorem 6.3 (b) implies that for any  $u \in M(S)$ ,  $n \ge 1$  and  $\varepsilon > 0$ , there exists a  $\pi$ -generated  $\pi'$  such that  $I_n(\pi', u) \ge U_n^n u - \varepsilon$ .

THEOREM 6.4. Under Assumptions (I) and (II),  $U_{\pi}$  is a contraction mapping on M(S) with the contraction coefficient K, i.e.  $\|U_{\pi}u-U_{\pi}v\| \le K\|u-v\|$ , so that, from the Banach fixed-point theorem,  $U_{\pi}$  has a unique fixed point  $u^*$ , and  $\|U_{\pi}^nu-u^*\| \le K^n\|u-u^*\|$  for all n.

PROOF.  $v \le u + ||u - v||$  yields

$$U_{\pi}v \leq U_{\pi}(u + ||u - v||) \leq U_{\pi}u + K||u - v||$$

using Theorem 6.2 (b). Interchange u and v to obtain

$$U_{\pi}u \leq U_{\pi}v + K||v-u||$$
,

completing the proof.

THEOREM 6.5. Let Assumptions (I), (II) be satisfied.

- (a) For any Markov policy  $\pi = \{f_1, f_2, \cdots\}$ , denoting by  $T_{f_n}$  associated with  $f_n$  and by  $U_{\pi} = \sup_{n \geq 1} T_{f_n}$ , the operator associated with  $\pi$ , the fixed point  $u^*$  of  $U_{\pi}$  is the optimal return among  $\pi$ -generated policies:  $I(\pi') \leq u^*$  for every  $\pi$ -generated  $\pi'$ , and for every  $\varepsilon > 0$  there is a  $\pi$ -generated f such that  $I(f^{(\infty)}) \geq u^* \varepsilon$ . Any f with  $T_f u^* \geq u^* \varepsilon(1 K)$  satisfies this inequality. Furthermore, any f with  $||T_f u^* u^*|| \leq \varepsilon(1 K)$  satisfies  $||I(f^{(\infty)}) u^*|| \leq \varepsilon$ .
  - (b) For any  $p \in P(S)$ ,  $\varepsilon > 0$ , there exists a  $(p, \varepsilon)$ -optimal policy which is stationary.
- (c) For any  $\varepsilon \ge 0$ , if there is an  $\varepsilon$ -optimal  $\pi^* = {\pi_1, \pi_2, \cdots}$ , there exists an  $\varepsilon/(1-K)$ -optimal policy which is stationary.
- (d) Denote for each  $a \in A$  by  $T_a$  the operator associated with  $f \equiv a$ . Any u with  $T_a u \leq u$  for all a is an upper bounded on reward:  $I(\pi) \leq u$  for all  $\pi$ .
- (e) If for every  $\varepsilon > 0$  there exists an  $\varepsilon$ -optimal policy, then the optimal return  $u^*$  is a Baire function and it satisfies the optimality equation  $u^* = \sup T_a u^*$ .
- (f) A policy  $\pi$  is optimal if and only if its return  $I(\pi)$  satisfies the optimality equation.

PROOF. (a) For any  $\pi$ -generated  $\pi' = \{g_1, g_2, \cdots\}$ , we have  $I(\pi') = T_{g_1}T_{g_2} \cdots T_{g_n}u_n$ , where  $u_n = I(g_{n+1}, g_{n+2}, \cdots)$ . Since each  $T_{g_i}$  is a contraction operator with contraction coefficient K,

$$\begin{split} \|T_{g_1}T_{g_2} \cdots T_{g_n}u_n - T_{g_1}T_{g_2} \cdots T_{g_n}u^*\| & \leq K^n \|u_n - u^*\| \\ & \leq K^n (\|g\| + \|u^*\|) \,. \end{split}$$

Thus  $T_{g_1}\cdots T_{g_n}u^*\to I(\pi')$  as  $n\to\infty$ . But  $T_{g_1}\cdots T_{g_n}u^*\le U_\pi^nu^*=u^*$ , so that  $I(\pi')\le u^*$ . From Theorem 6.3 (a), there is a  $\pi$ -generated f for which  $T_fu^*\ge U_\pi u^*-\varepsilon'=u^*-\varepsilon'$ , where  $\varepsilon'=\varepsilon(1-K)$ . We can verify inductively that

$$T_f^n u^* \ge u^* - \varepsilon' (1 + K + \cdots + K^{n-1})$$
 for all  $n \ge 1$ .

Since  $T_f^n u^* \to I(f^{(\infty)})$ , we conclude that

$$I(f^{(\infty)}) \ge u^* - \lceil \varepsilon'/(1-K) \rceil = u^* - \varepsilon$$
.

Last two inequalities are immediate from Theorem 6.3 (a).

(b) This is immediate from Corollary to Theorem 5.1 and above (a).

(c) For any  $\pi^* = \{\pi_1, \pi_2, \cdots\}$ ,  $I(\pi^*) = e_{\pi^*}g = \pi_1 qg(s, a, s_2, E^{1\pi^*}g)$ , where  $E^{1\pi^*}g = E^{1\pi^*}g(s, a, s_2) = I(\pi_{s,a})(s_2)$  and  $\pi_{s,a}$  denotes the policy which  $\pi^*$  specifies starting with the second stage, when the first state and action are s, a respectively, i.e.  $\pi_{s,a} = \{\pi'_1, \pi'_2, \cdots\}$  where

$$\pi'_n(\cdot | s_1, a_1, \dots, s_n) = \pi_{n+1}(\cdot | s, a, s_1, a_1, \dots, s_n)$$
.

If  $\pi^*$  is  $\varepsilon$ -optimal,  $E^{1\pi^*}g(s, a, s_2) = I(\pi_{s,a})(s_2) \leq I(\pi^*)(s_2) + \varepsilon$  for all  $s_2 \in S$ , so that, by monotonicity of g

$$\begin{split} I(\pi^*) &= \pi_1 q g(s_1, \ a_1, \ s_2, \ E^{1\pi} g) \\ &\leq \pi_1 q g(s_1, \ a_1, \ s_2, \ I(\pi^*)(s_2) + \varepsilon) \\ &= \pi_1 h \ , \qquad \text{say} \ . \end{split}$$

From the Lemma 2.1, there exists an f for which  $\pi_1 h \leq fh$  for all s, namely

$$\pi_1 qg(s_1, a_1, s_2, I(\pi^*)(s_2) + \varepsilon) \leq fqg(s_1, a_1, s_2, I(\pi^*)(s_2) + \varepsilon)$$

$$= T_f(I(\pi^*) + \varepsilon).$$

Then,

$$I(\pi^*) \leq T_f(I(\pi^*) + \varepsilon) \leq T_fI(\pi^*) + K\varepsilon$$
.

Again,

$$T_f(I(\pi^*)) \leq T_f(T_f(I(\pi^*) + K\varepsilon)) \leq T_f^2 I(\pi^*) + K^2 \varepsilon$$
.

By induction on n we obtain

$$I(\pi^*) \leq T_f^n I(\pi^*) + K \varepsilon + K^2 \varepsilon + \cdots + K^n \varepsilon$$
.

Letting  $n \rightarrow \infty$  yields

$$I(\pi^*) \leq I(f^{(\infty)}) + K\varepsilon/(1-K)$$
.

Since  $\pi^*$  is  $\varepsilon$ -optimal,  $f^{(\infty)}$  is  $\varepsilon/(1-K)$ -optimal.

- (d) The proof of (d) is similar to that of Theorem 6 (d) in [1], and is omitted.
- (e) From (c), the hypothesis implies that there exists a 1/n-optimal stationary policy  $f_n^{(\infty)}$ , say. With  $\pi=\{f_1,f_2,\cdots\}$ , the fixed  $u^*$  of the  $U_\pi$ , is, from (a), the optimal reward among  $\pi$ -generated policies. In particular  $u^* \geq I(f_n^{(\infty)})$ , so that  $u^* \geq I(\pi)$  for all  $\pi$ , and  $u^*$  is the optimal return. We have  $\sup_{a \in A} T_a u^* \geq U_\pi u^* = u^*$ . On the other hand, for any  $a \in A$ ,

$$\begin{split} T_a u^* & \leq T_a(I(f_n^{\scriptscriptstyle(\infty)}) + (1/n)) \leq T_a(I(f_n^{\scriptscriptstyle(\infty)}) + \frac{K}{n}) \\ & = I(a, f_n^{\scriptscriptstyle(\infty)}) + \frac{K}{n} \leq u^* + \frac{K}{n} \;, \end{split}$$

where  $(a, f_n^{(\infty)})$  is the Markov policy  $\{g, f_n, f_n, \dots\}$  with  $g \equiv a$ . Letting  $n \to \infty$  yields  $T_a u^* \leq u^*$ . Thus  $u^*$  satisfies the optimality equation.

(f) The proof is same as in Theorem 6 (f) of [1].

Finally we give a theorem in the case when A is countable or finite, without stating the proof.

Theorem 6.6. Let Assumptions (I) and (II) be satisfied. (i) If A is countable, then for any  $\varepsilon > 0$  there exists an  $\varepsilon$ -optimal policy which is stationary. (ii) If A is finite, then there exists an optimal policy which is stationary.

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